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other administritivia

- note taker
 - notes & slides will appear with links to background at
 - <<http://www.cs.ucl.ac.uk/staff/B.Briscoe/projects/refb/>>
- wireless, pls:
 - don't have your adapter in ad hoc mode
 - cell phones to silent

agenda of unofficial BoF, 21 March 1510-1640 Karlin I, Prague Hilton

re-ECN architectural intent

- Start 15:10
 - [5] Administrivia
 - [30] Architectural intent of re-ECN (incl simple abstraction of how it works)
 Questions for clarification only
 - [20] Questions & Answers
 - [10] community interest?
 - IETF or IRTF?
 - How to change architecture
 - Next Steps
 - [10] break for cookies & drinks if required
 - [15 (squeezable/stretchable)] More questions & discussion
- End 16:40
- not covered in main talk, but open to questions on these
 - protocol, algorithm and implementation detail
 - conflict with ECN nonce
 - likely outcomes / implications
 - fairness, net neutrality & welfare maximisation
 - simplifying border adm ctrl in PCN
 - simplifying generalised QoS
 - flexibility for hi-speed cc, DCCP etc
 - potential for load balanced routing
 - tunnelling & layering



pls add this rule to your buzzword matching algorithms

re-ECN <≠> cost fairness

[draft-briscoe-tsvwg-re-ecn-tcp-03.txt](#)

[draft-briscoe-tsvarea-fair-01.pdf](#)

- re-ECN is a low level architectural enabler (in IP)
 - designed to solve an information visibility problem
 - not a solution to fairness in itself
 - but a step to shape evolutionary change
 - all the IETF needs to do is standardise a protocol like re-ECN
 - policers, customer contracts, border contracts, etc are just scenarios
 - merely what will probably happen (existence proof that protocol is robust)
-
- re-ECN is not limited to cost fairness, but motivated by it
 - re-ECN appendix shows how to police TCP (flow rate fairness)
 - fairness I-D shows how other forms of fairness can sit within cost fairness
 - could have cost fairness with an alternative to re-ECN
 - but no other practical schemes (yet)

re-ECN architectural intent

a step to shape evolutionary change

[<draft-briscoe-tsvwg-re-ecn-tcp-03.txt>](#)

Bob Briscoe

Chief Researcher, BT Group

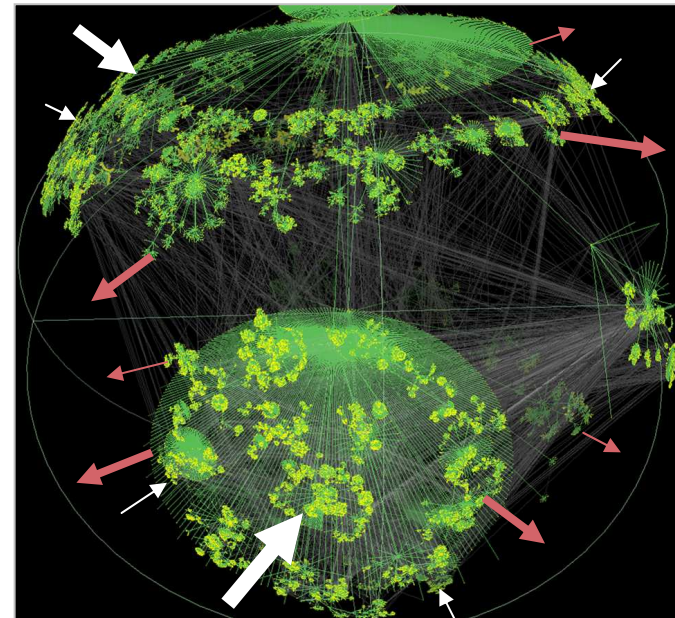
unofficial Birds of a Feather at IETF-68

Mar 2007



known problem since early days

- how to share all the parts of a huge, multi-provider packet multiplexer between competing processes
- keeping one-way datagrams
- allowing for
 - self-interest & malice
 - of users and of providers
 - evolvability
 - of new rate dynamics from apps
 - of new business models
 - viability of supply chain
 - simplicity

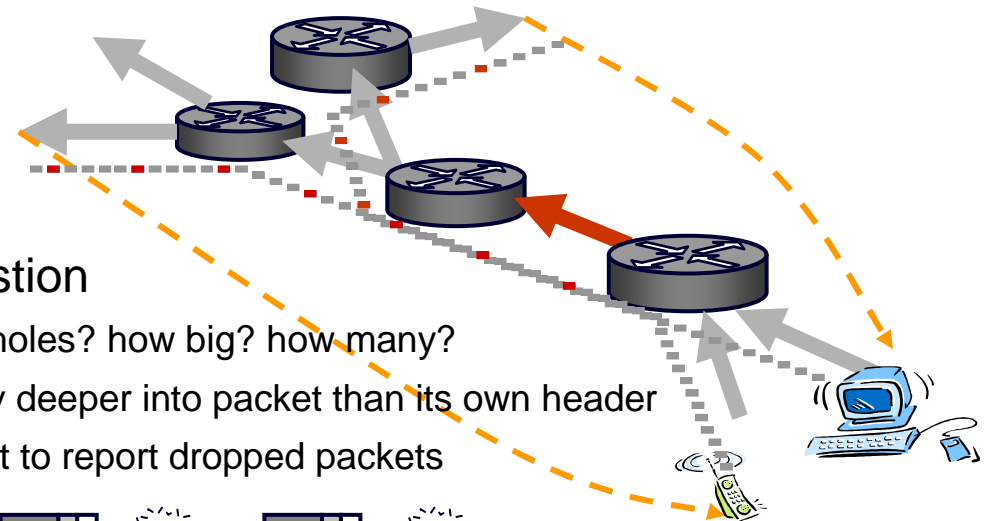


Internet topology visualization produced by Walrus
(Courtesy of Young Hyun, [CAIDA](#))

- if we do nothing
 - the few are ruining it for the many
 - massive capacity needed to keep interactive apps viable
 - poor incentives to invest in capacity
 - operators are kludging it with DPI
 - solely today's apps frozen into net
 - complex, ugly feature interactions

solution step #1: ECN

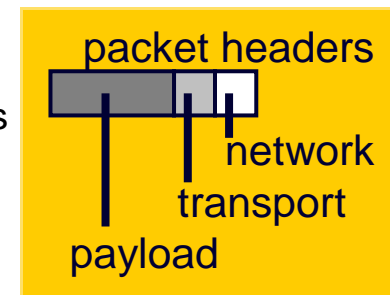
make congestion visible to network layer



- packet drop rate is a measure of congestion
 - but how does network at receiver measure holes? how big? how many?
 - can't presume network operator allowed any deeper into packet than its own header
 - not in other networks' (or endpoints') interest to report dropped packets



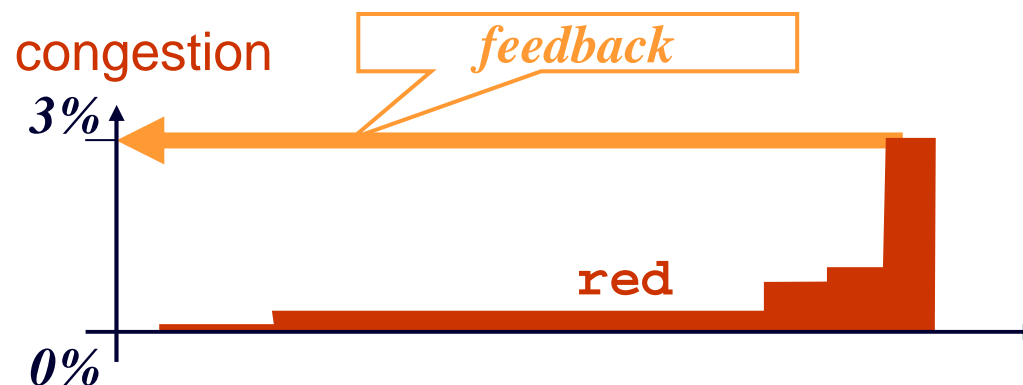
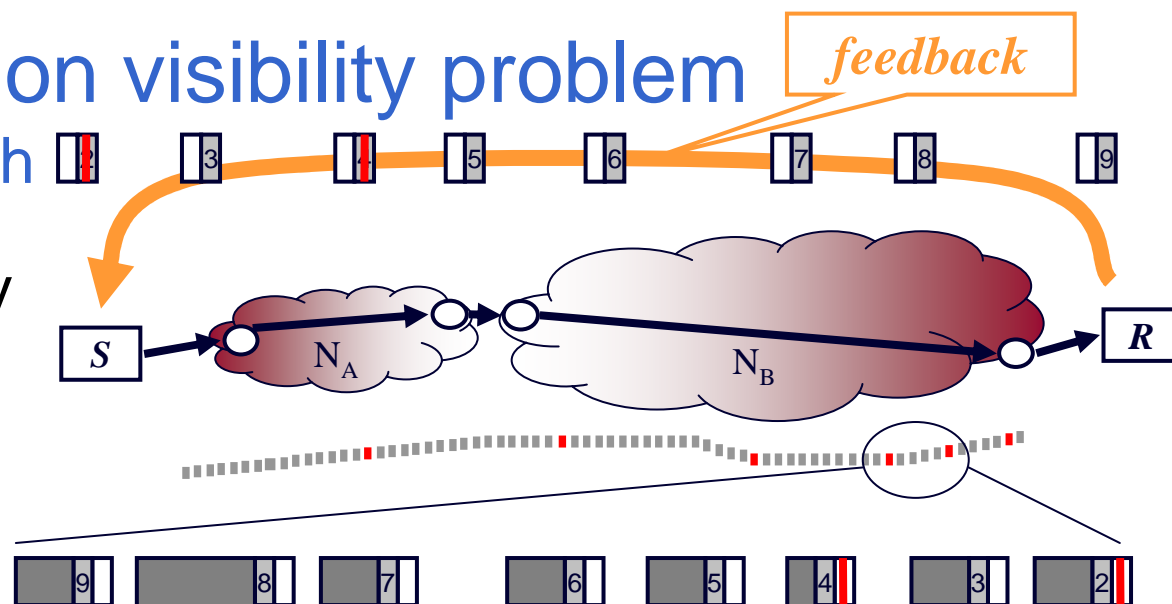
- solution: Explicit Congestion Notification (ECN)
 - mark packets as congestion **approaches** - to avoid drop
 - already standardised into IP (RFC3168 – 2001)
 - implemented by most router vendors – very lightweight mechanism
 - but rarely turned on by operators (yet) – mexican stand-off with OS vendors



new information visibility problem

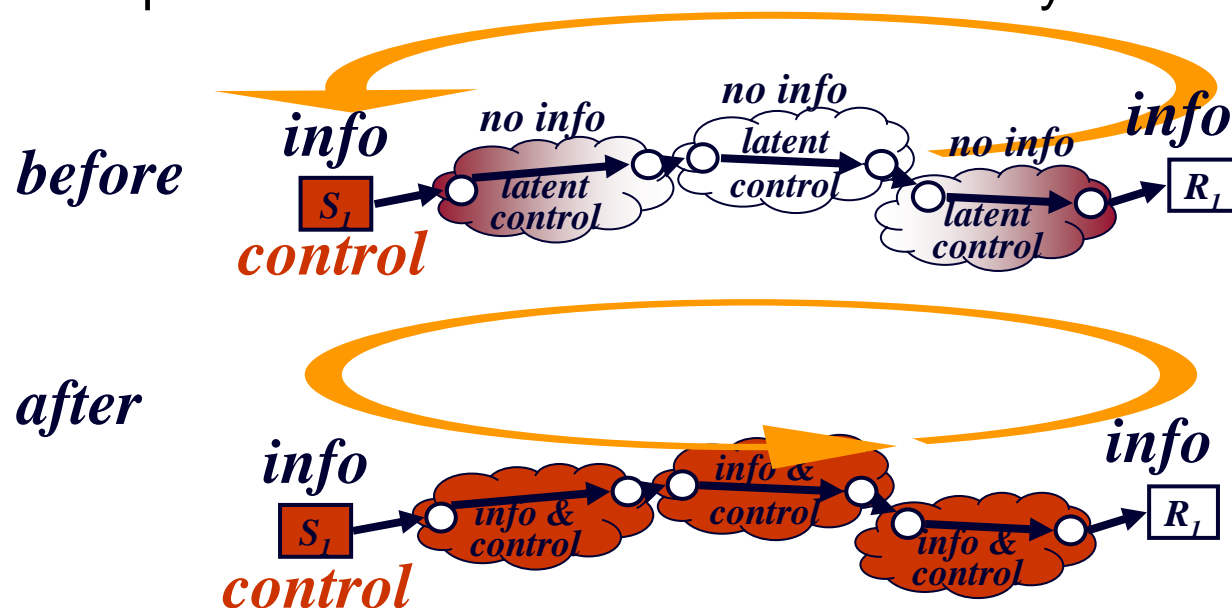
ECN is not enough

- path congestion only measurable at exit
- can't measure path congestion at entry
 - can't presume allowed deeper into feedback packets



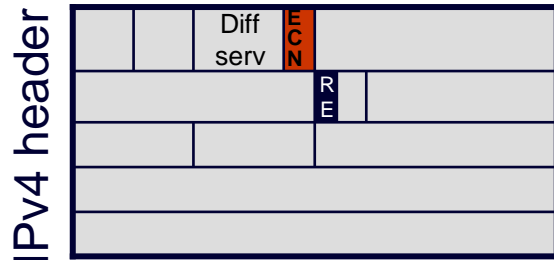
re-ECN in brief

- reinsert feedback
- packets arrive at each router predicting downstream path
- incremental deployment + upgrade incentive knob
- hangs new capabilities on ECN deployment, not just performance
- a simple idea for the Internet's accountability architecture



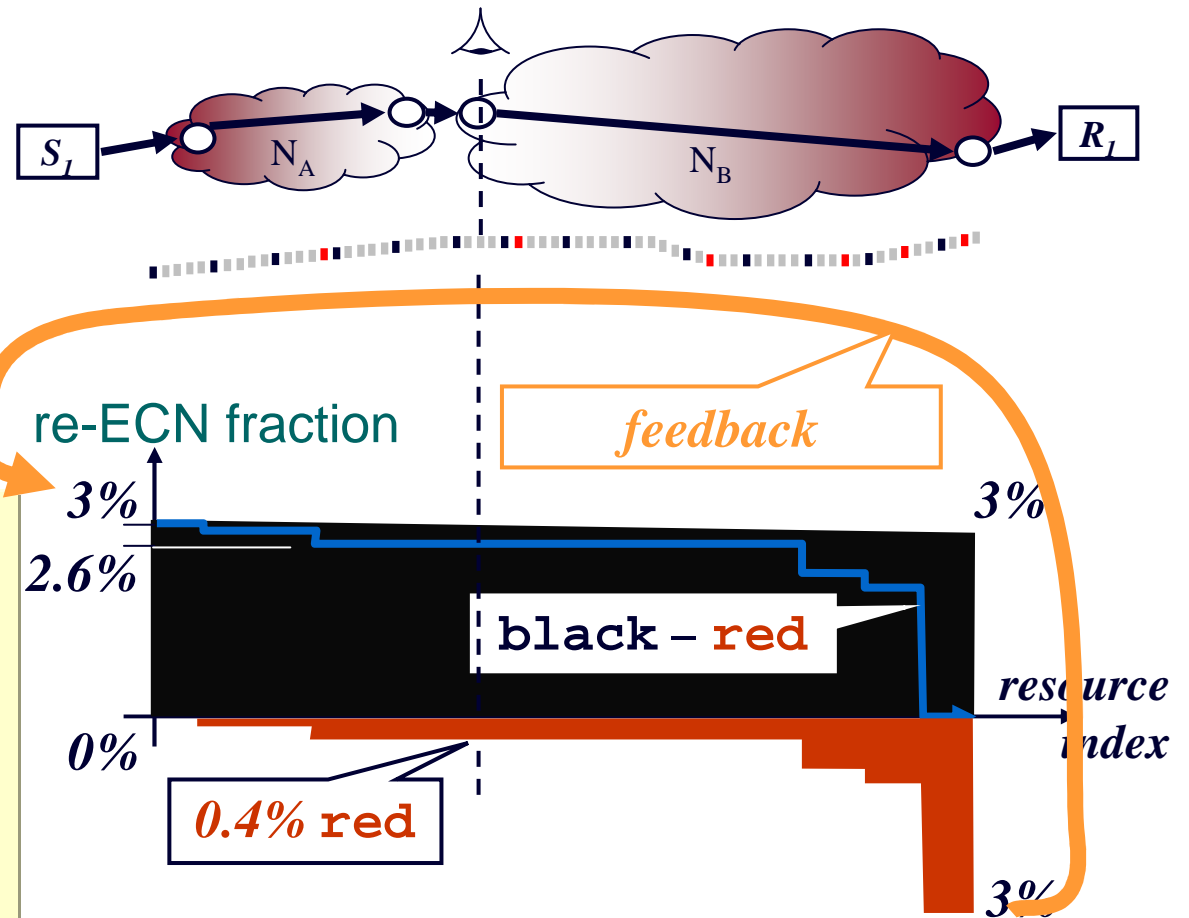
measurable downstream congestion

solution step #2



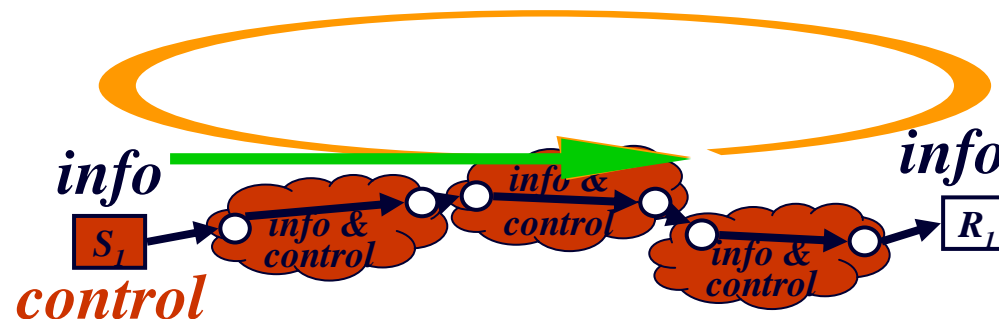
re-feedback

- sender re-inserts feedback by marking packets **black**
- at any point on path, diff betw fractions of **black** & **red** bytes is downstream congestion
- ECN routers unchanged
- **black** marking e2e but visible at net layer for accountability



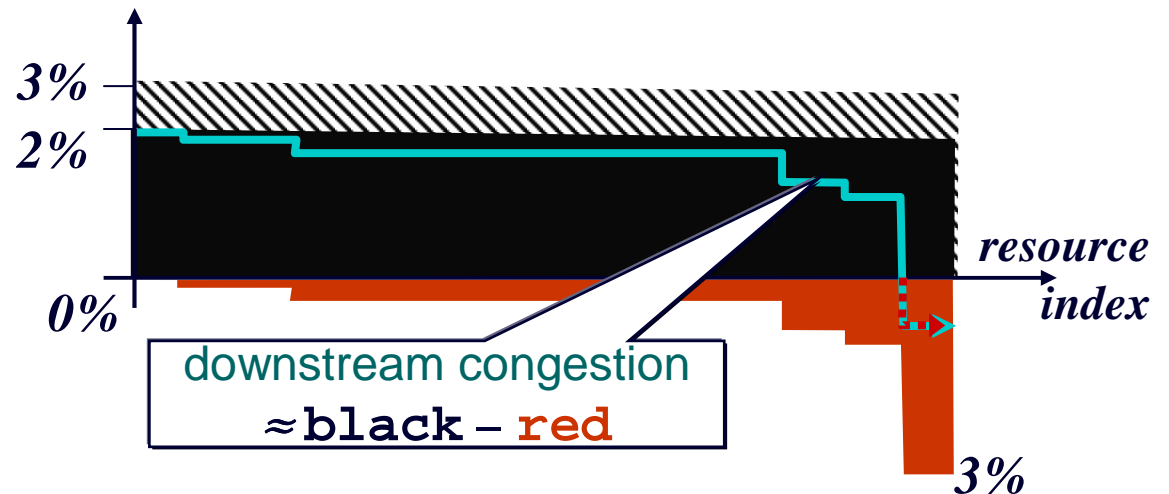
flow bootstrap 'pre-feedback'

- at least one **green** packet(s) at start of flow or after >1sec idle
 - means “feedback not established”
 - ‘credit’ for safety due to lack of feedback
 - a **green** byte is ‘worth’ same as a **black** byte
- lots of powerful uses for a different colour from black
 - distinguishes conservatism from expected congestion based on experience
 - ability to vary the expected cost of jump-starting (research needed)
 - gives deterministic flow state mgmt (policers, droppers, firewalls, servers)

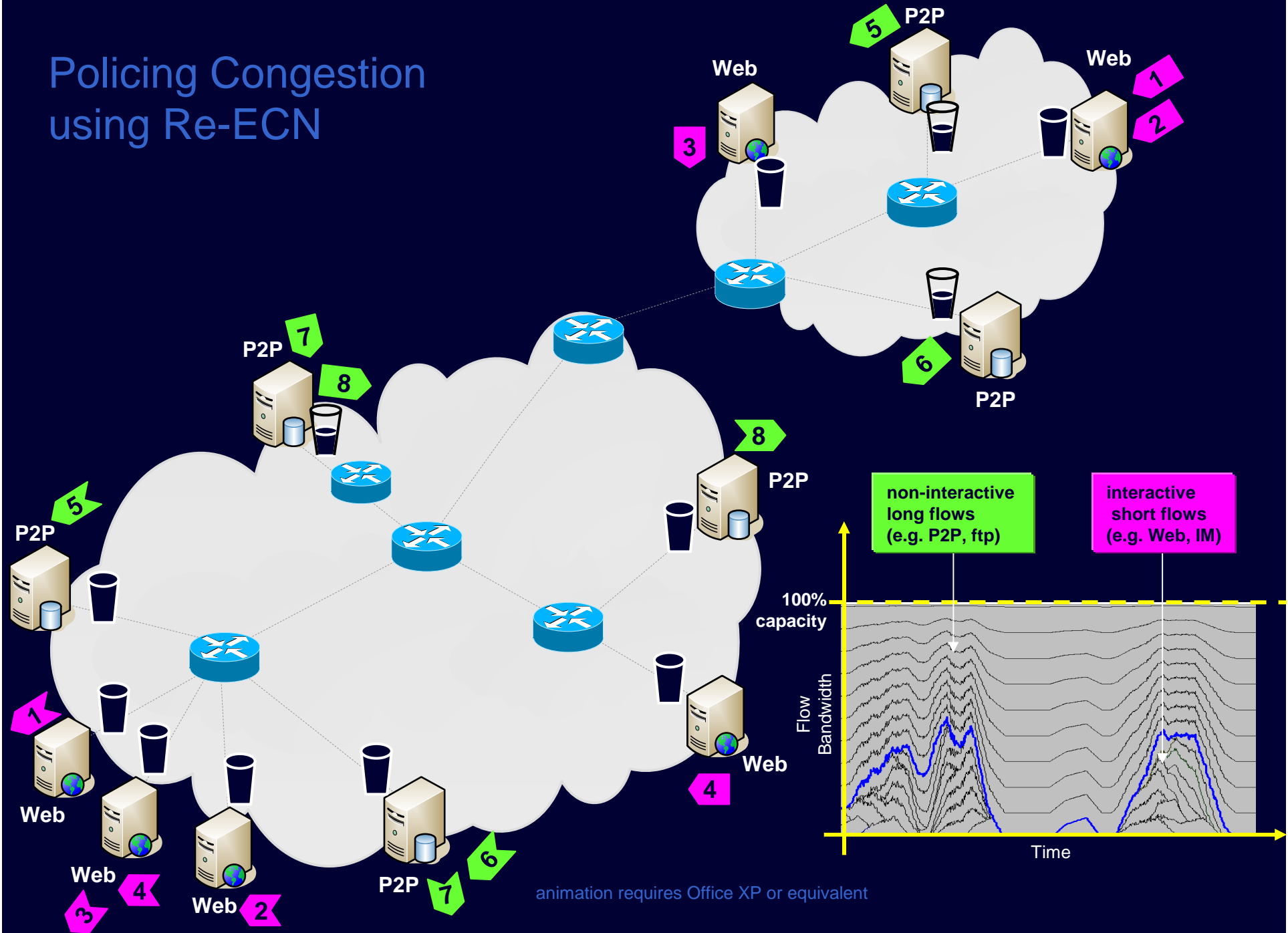


proposed re-ECN service model

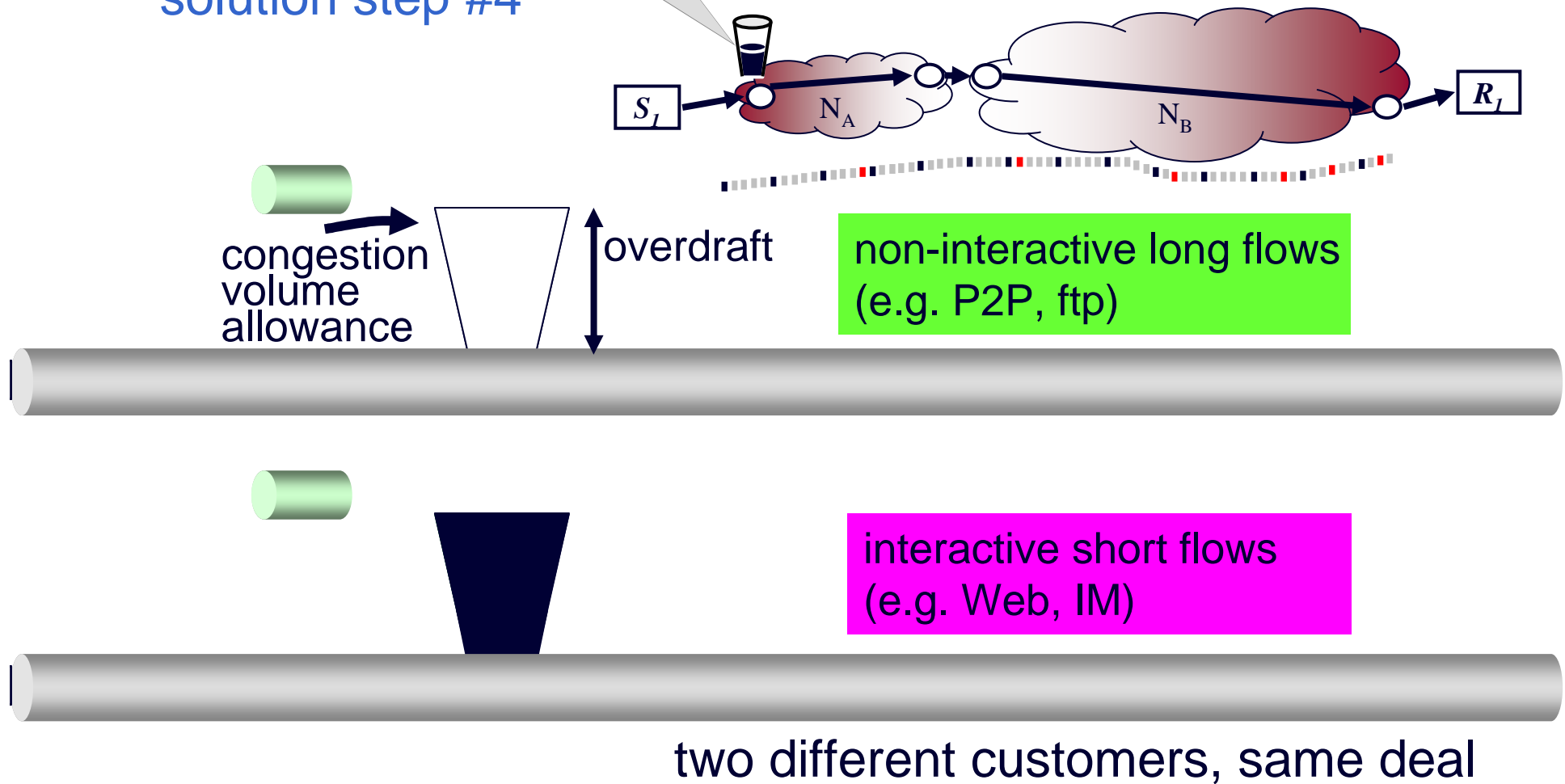
- to encourage sender (or proxy) to indicate sufficient expected congestion...
- Internet won't try to deliver packet flows beyond the point where more congestion has been experienced than expected
 - if sender wants to communicate, has to reveal expected congestion
 - even if sender not trying to communicate (e.g. DoS) packets can be dropped rather than enqueued before they add to congestion



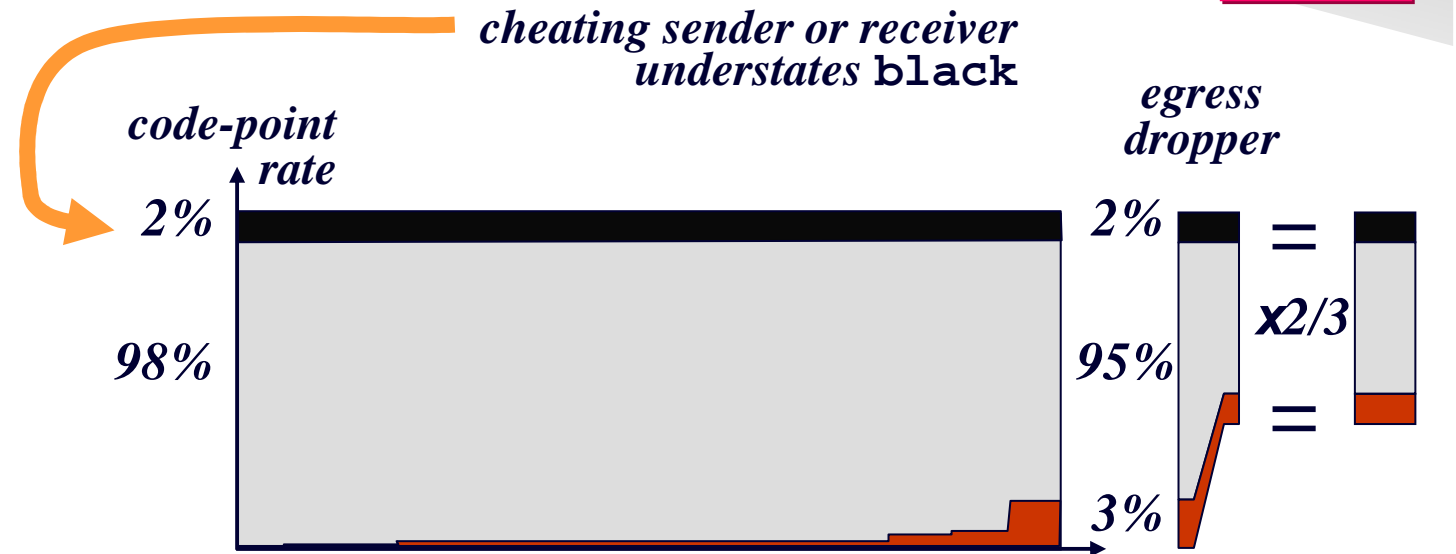
Policing Congestion using Re-ECN



congestion policer – one example: per-user policer
solution step #4



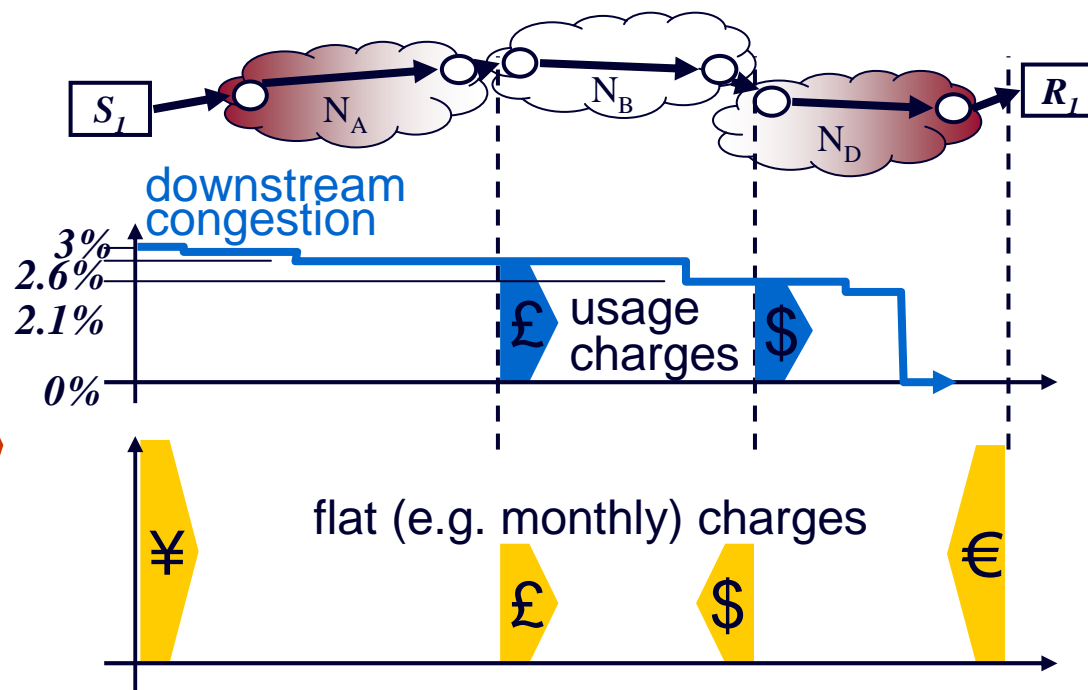
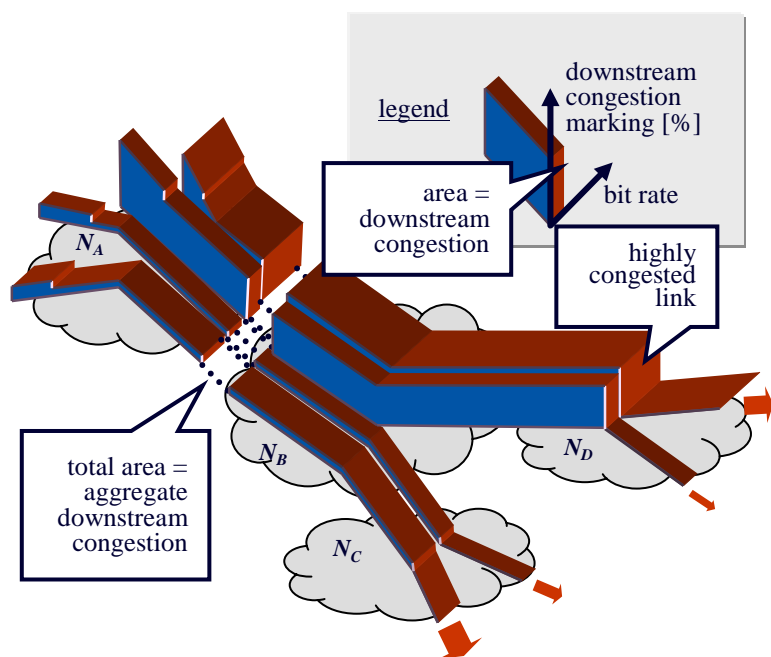
egress dropper (sketch)



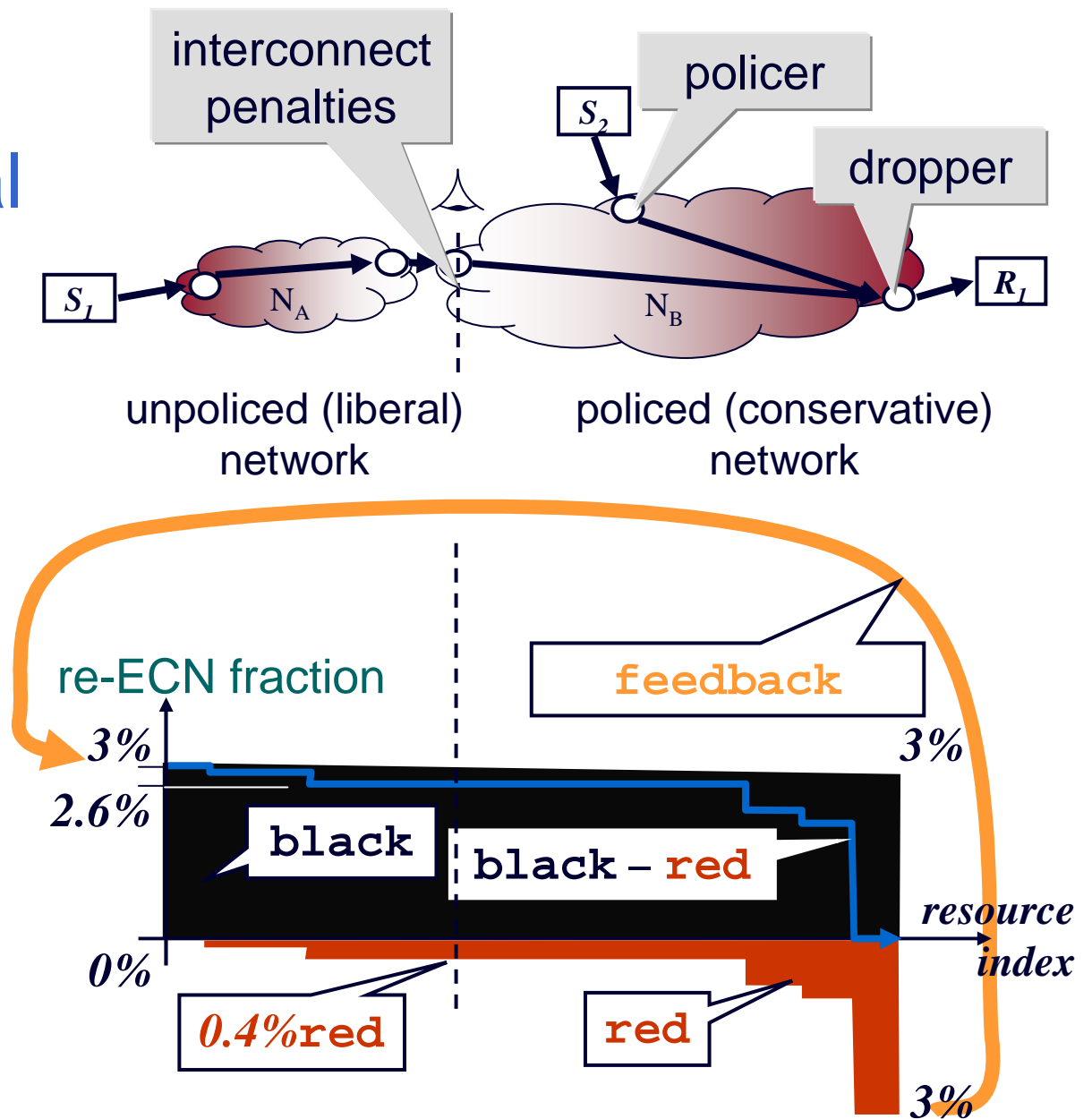
- drop enough traffic to make fraction of **red** = **black**
 - understatement allows gain through policer, but dropper always fully cancels it out
 - goodput best if rcvr & sender honest about feedback & re-feedback
- understate congestion to attack routers?
 - given overloaded routers, honest senders will be sending nearly all **black**
 - overloaded routers preferentially drop **grey** and **red** (next slide)
- important principle: attack traffic does no harm until it congests a router
 - re-ECN drops attack at first congested router (no push-back, no new attack vector)

inter-domain accountability for congestion

- metric for inter-domain SLAs or usage charges
 - N_B applies penalty to N_A in proportion to bulk volume of **black** less bulk volume of **red** over, say, a month
 - could be tiered penalties, directly proportionate usage charge, etc.
 - flows de-aggregate precisely to responsible networks



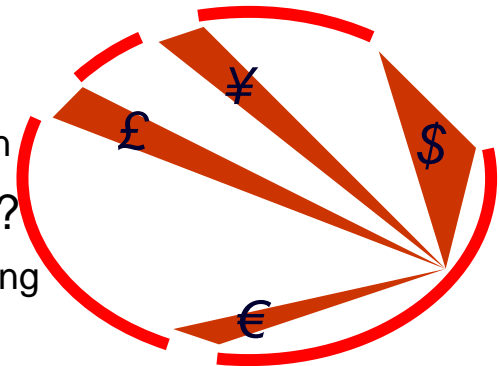
re-ECN partial deployment



deployment incentives

bootstrap then chain reaction

- deployment effectively involves architectural change
 1. (minor) change to sender's Internet stack
 2. network deploys edge/border incentive functions
 - breaking the stand-off between 1 & 2 requires strong incentives
- re-feedback solves ISPs' main cost control problem
 - third party services competing with ISP pay below network cost
 - ISP has to compete *while* paying balance of competitor's costs
 - hits big fear button and big greed button
 - but keeps moral high ground
 - net neutral: managing congestion not app discrimination
- first movers: vertically integrated cellular operators?
 - 3GPP devices leak deployment to other networks by roaming
- 2nd movers (NGNs?) continue chain reaction
 - adopters' incoming border charges focus on non-adopters



outstanding issues

- technical

- ✕ a lot more verification of all the claims to do
- ✕ community found a few nasty vulnerabilities over last two years
 - ✓ fixed (added minor complexity in only one case)
- ✕ connection spoofing attack still outstanding
 - ✓ possible solution recently brainstormed

- religious

- ✕ underlying problem has been dogma that equal flow rates are fair
 - ✓ groundswell change in community thinking since mid Oct'06
- ✕ dismantling a religion not so easy

- community

- ✕ a lot of passive support but consensus needs a lot more active interest

a change to IP needs to be 'owned' by Internet community
please take it, break it, analyse it, re-design it, work out implications

conclusions

- resolution of tensions in fairness / net neutrality debate
 - freedom to use the Internet, until you congest freedom of others
 - proportionate restriction of freedom during congestion
- an architectural change with grand implications
 - simple management and control of fairness & QoS
 - naturally mitigates DDoS
 - generates correct capacity investment incentives and signals
- but conceptually simple and trivial to implement
- strong deployment incentives
 - bootstrap and onward chain reaction
- where's the catch?
 - invite you to analyse it, break it, re-design it



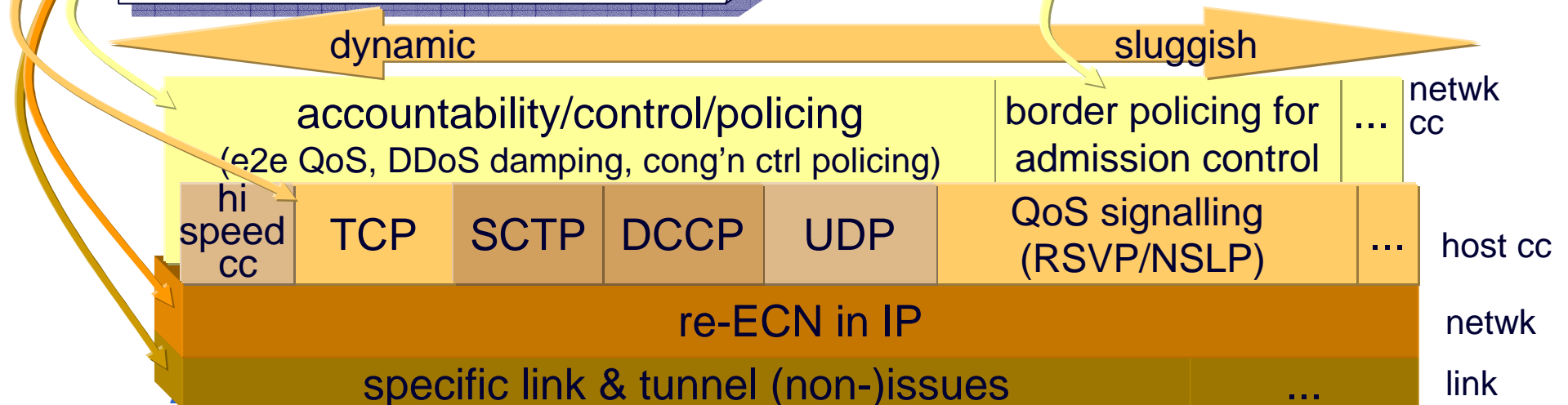
Internet draft roadmap

Re-ECN: Adding Accountability for Causing Congestion to TCP/IP
[draft-briscoe-tsvwg-re-ecn-tcp-03](#)

§3: overview in TCP/IP *intent*
§4: in TCP & other transports *stds*
§5: in IP (v4 & v6)
§6: accountability apps *inform'l*

•more papers (PCN, QoS, DDoS etc):
<<http://www.cs.ucl.ac.uk/staff/B.Briscoe/projects/refb/>>

Emulating Border Flow Policing using Re-ECN on Bulk Data
[draft-briscoe-tsvwg-re-ecn-border-cheat-02](#)
intent: informational



re-ECN architectural intent

<<http://www.cs.ucl.ac.uk/staff/B.Briscoe/projects/refb/>>

spare slides

more motivating problems

more architectural motivation

- (non)issues with layering & tunnelling
- bottleneck policing harmful
- independence from identifiers

mechanism

- IPv4 & v6 wire protocol
- drop preference semantics
- conflict with the ECN nonce

uses

- simplifying generalised QoS
- flexibility for hi-speed cc, DCCP etc
- adding re-ECN to various transports: TCP, SCTP, DCCP, PCN, UDP
- DDoS mitigation
- potential for load balanced routing

incentives and security

- attacks on re-ECN and fixes
- IPR

Q&A

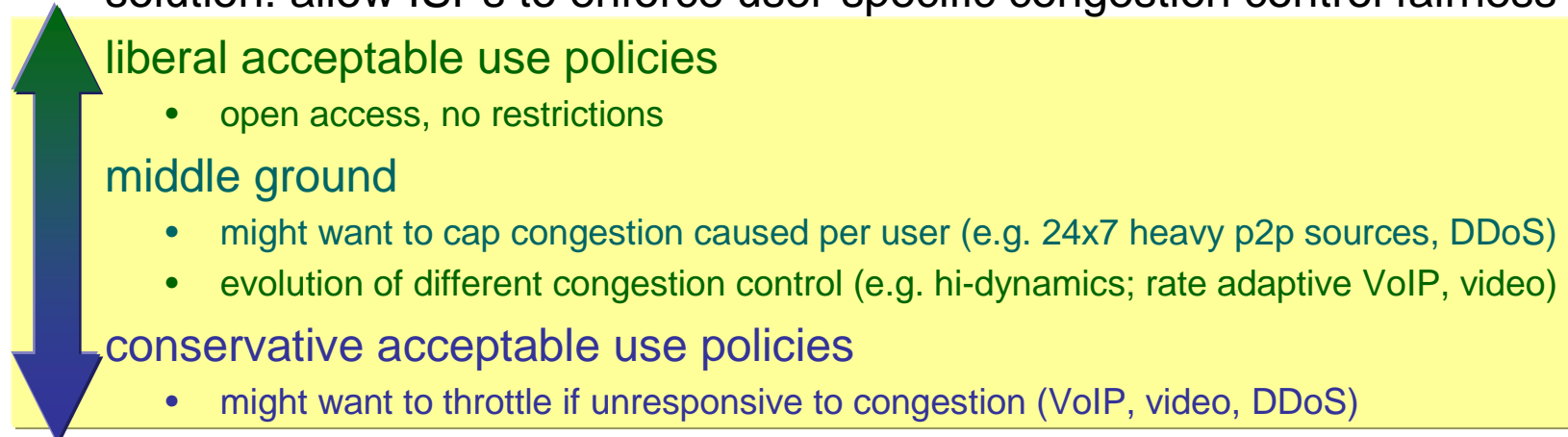


next steps

- build community
- simulations, implementation continues
- Official IETF BoF?
- IRTF Internet Congestion Control research group?

designed for tussle

- current Internet gives freedom but no fairness
 - the more you take, the more you get; the more polite you are, the less you get
 - but we don't want to lose freedom by enforcing fairness
- solution: allow ISPs to enforce user-specific congestion control fairness



- engineers shouldn't pre-judge answer to these socio-economic issues
 - Internet needs all these answers – balance to be determined by natural selection
 - 'do-nothing' doesn't maintain liberal status quo, we just get more middlebox kludges
- re-ECN at network layer: goals
 - just enough support for conservative policies without breaking 'net neutrality'
 - nets that allow their users to cause congestion in other nets can be held accountable

designed for tussle

Internet needs all these answers – market selection finds balance

demand side – freedom to degrade others

- the Internet is all about the freedom to get what I want (within my line rate)

limited by how much I impinge on the freedom of others

- enforceable congestion control

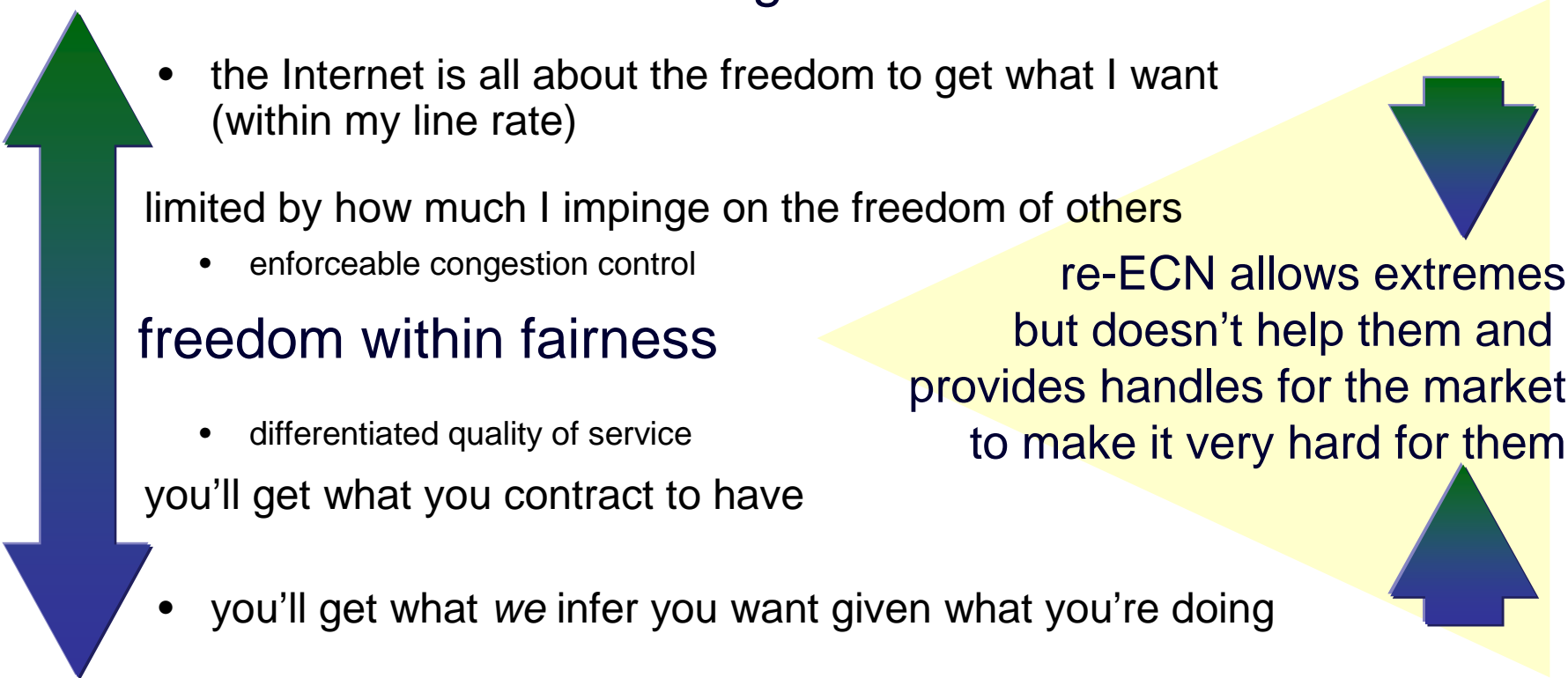
freedom within fairness

- differentiated quality of service

you'll get what you contract to have

- you'll get what we infer you want given what you're doing

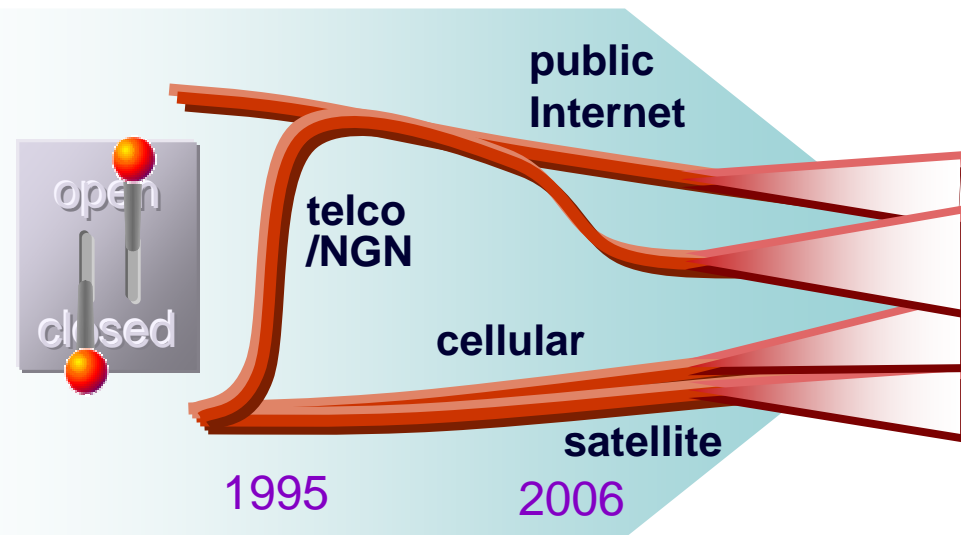
supply side – freedom to degrade competitors



re-ECN allows extremes
but doesn't help them and
provides handles for the market
to make it very hard for them

summary

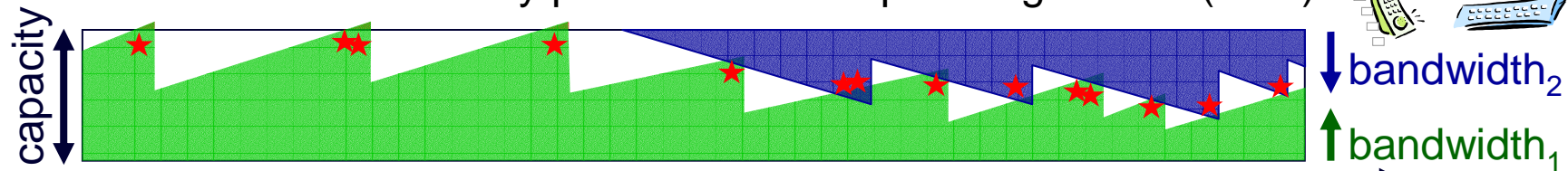
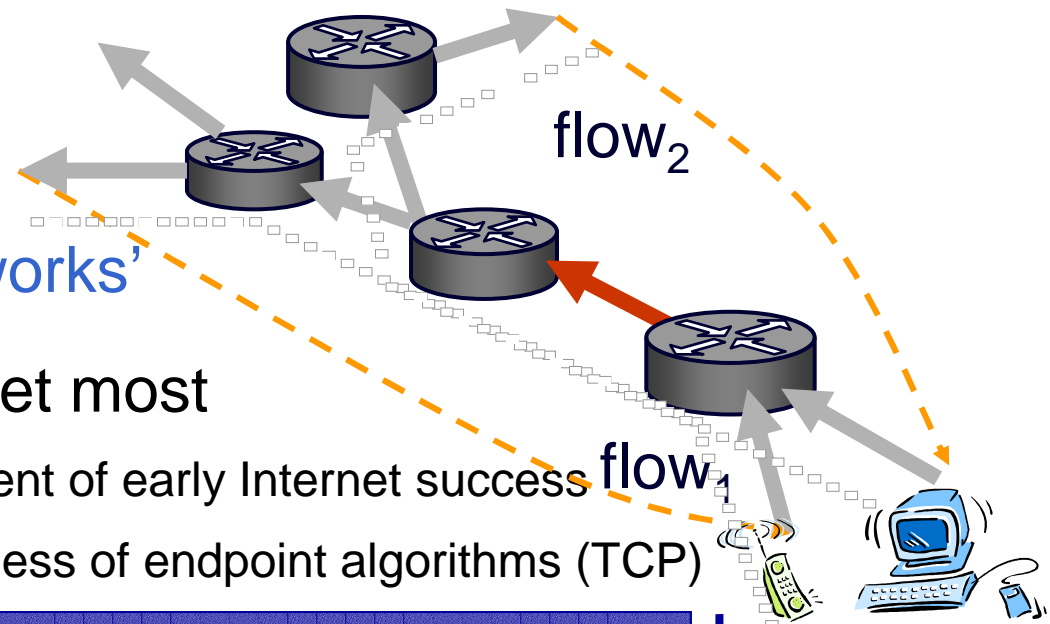
- Internet needs to be able to discriminate
 - against bits limiting the freedom of others – *bits* causing congestion
 - then wouldn't need to discriminate against *apps* causing congestion
- operators can choose not to limit their users' freedoms
 - but they take responsibility for congestion their users cause in other nets
- if operators do discriminate against apps
 - customers need enough choices to be able to switch operators
 - or apps can often obfuscate themselves anyway
- these economic effects require change to the Internet Protocol
 - making IP more suitable as the basis of a converged architecture



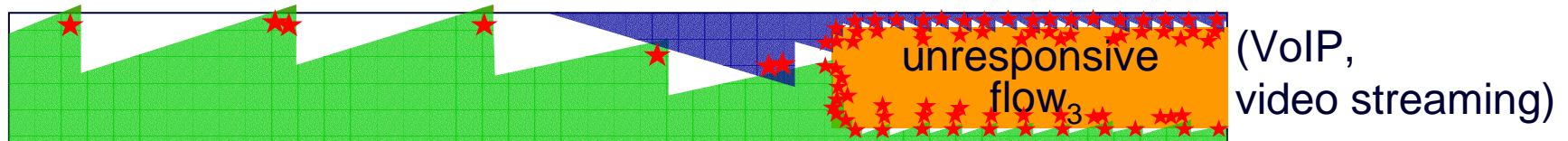
freedom

how Internet sharing 'works'

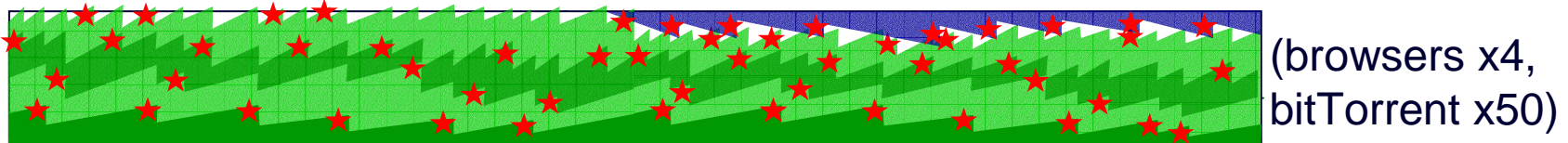
- those who push most, get most
 - restraint: the other ingredient of early Internet success
 - reliant on voluntary politeness of endpoint algorithms (TCP)



- a game of chicken – taking all and holding your ground pays



- or starting more 'TCP-fair' flows than anyone else



- or for much longer than anyone else (p2p file-sharing)

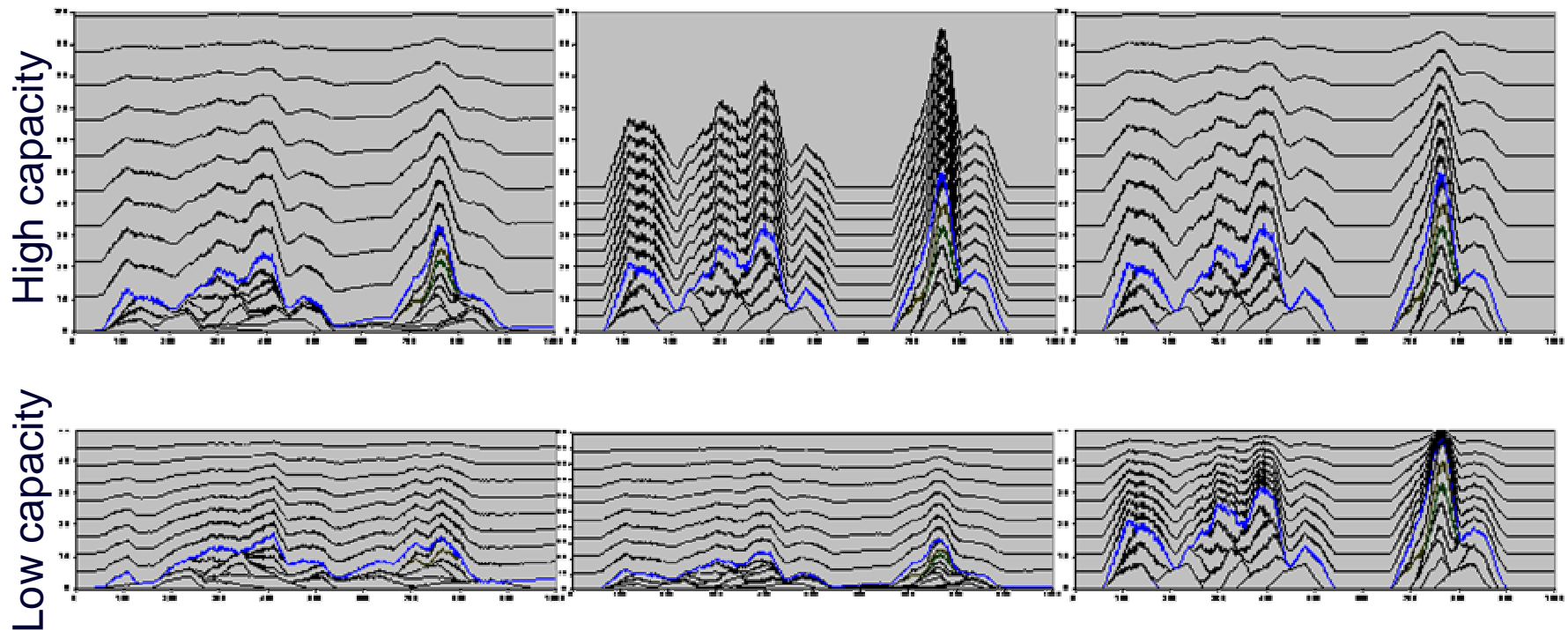
congestion cap auto-adjusts

volume cap always a hard compromise

No cap or loose volume cap

Tight volume cap

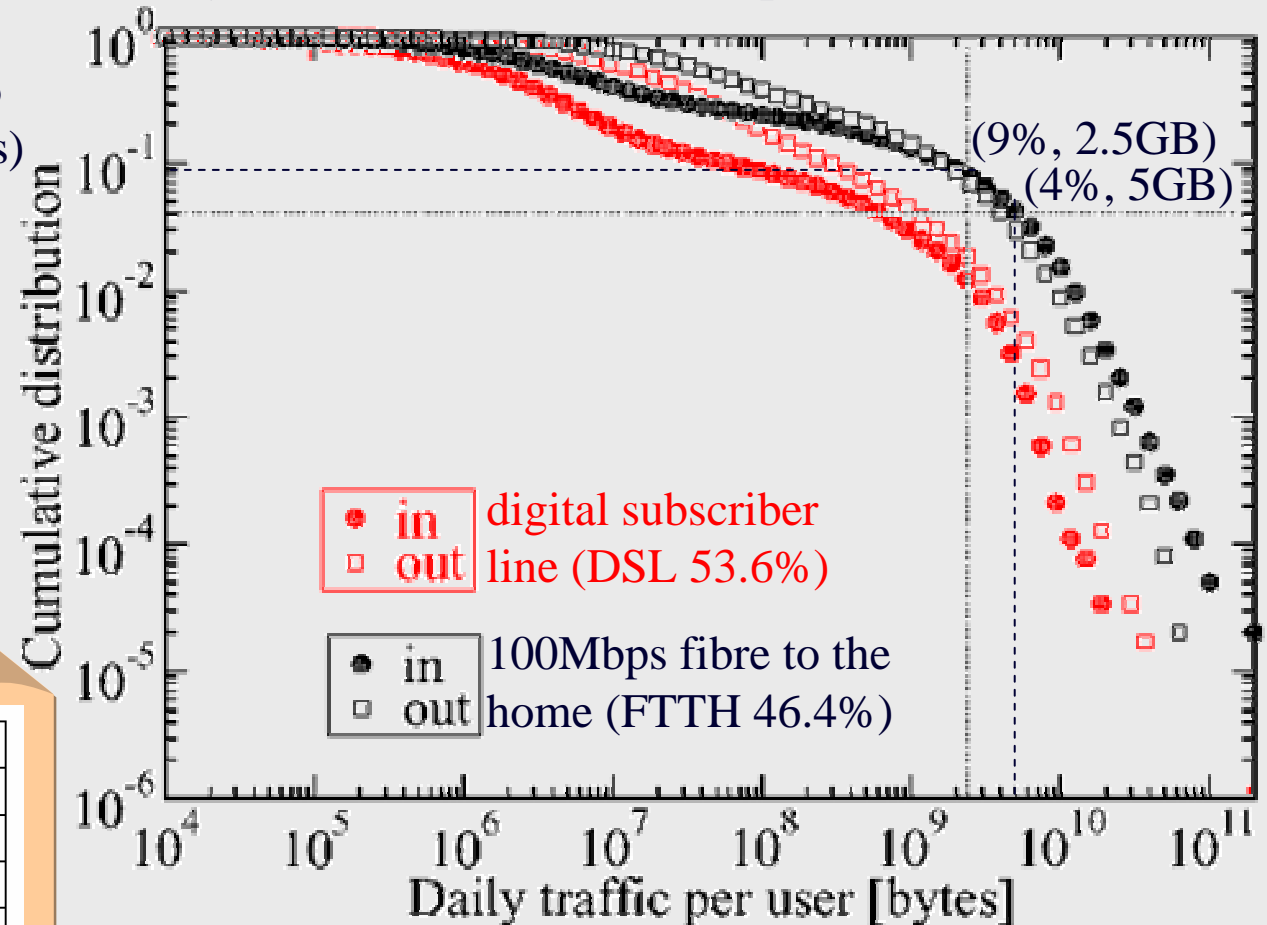
Congestion allowance



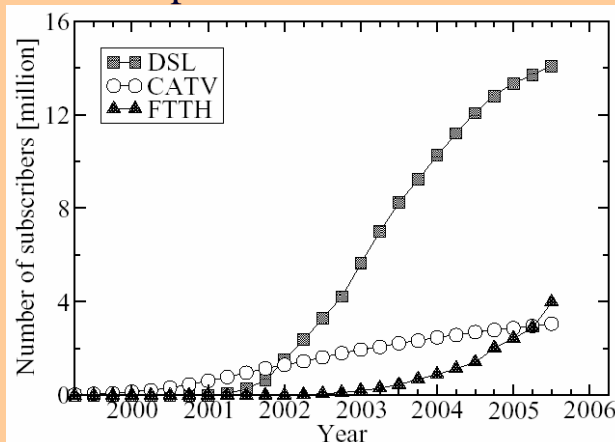
capacity growth will prevent congestion?

Distribution of customers' daily traffic into & out of a Japanese ISP (Feb 2005)

(5GB/day equivalent to
0.46Mbps if continuous)



Changing technology shares
of Japanese access market



Courtesy of Kenjiro Cho et al
The Impact and Implications of the Growth
in Residential User-to-User Traffic, SIGCOMM'06

hang on! solution step #0:

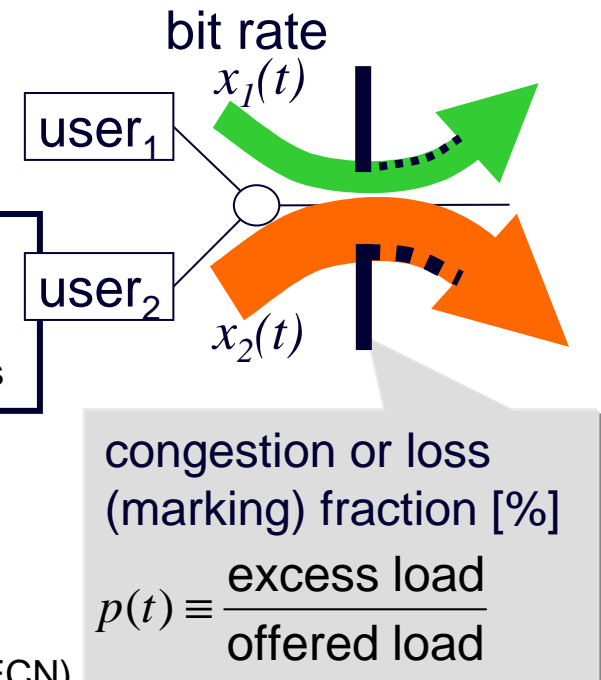
what's congestion got to do with the problem?

- can't solve a sharing problem without sharing costs
 - congestion is the cost of usage
 - cost of your behaviour on others

NOTE WELL

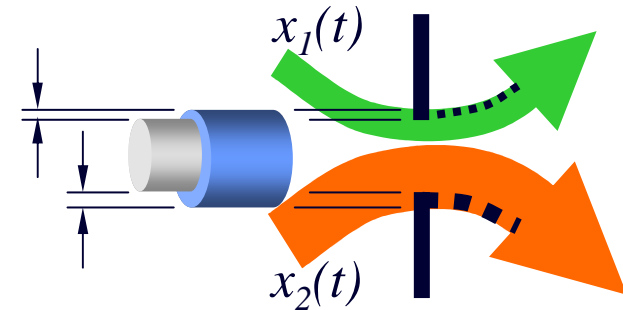
- IETF needs to provide the cost metric
- don't need metric for value – leave that for industry to guess

- it's not what you get
it's what you unsuccessfully tried to get
 - the bits each you contributed to excess load
 - loss/marking fraction $p(t)$ times your bit rate $x_i(t)$
 - only need dimensionless loss fraction $p(t)$ in wire protocol (ECN)
 - it subtly communicates your excess rate, because your own rate $x_i(t)$ is visible
- excess bits accumulate simply and correctly
 - over time, over flows and over network paths
 - congestion volume = bits of dropped/marked data you sent



calibrating 'cost to other users'

- a monetary value can be put on 'what you unsuccessfully tried to get'
 - the marginal cost of upgrading network equipment
 - so it wouldn't have marked the volume it did
 - so your behaviour wouldn't have affected others
- competitive market matches...
 - the cost of congestion volume
 - with the cost of alleviating it



*note: diagram is conceptual
congestion volume would be
accumulated over time*

*capital cost of equipment would
be depreciated over time*

- congestion volume is not an extra cost
 - part of the flat charge we already pay
 - but we can't measure who to blame for what
 - if we could, we **might** see pricing like this...

access link	congestion volume allow'ce	charge
100Mbps	50MB/month	€15/month
100Mbps	100MB/month	€20/month

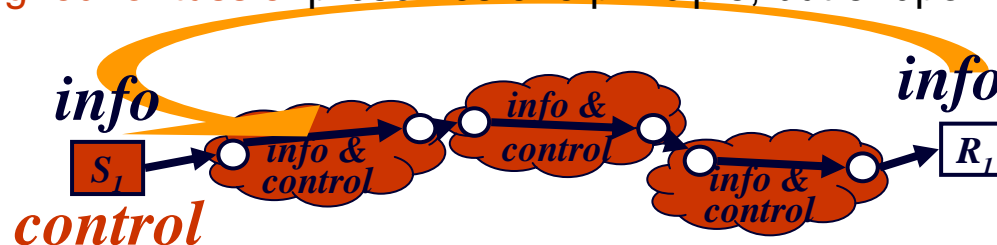
- NOTE WELL
 - IETF provides the metric
 - industry does the business models

re-feedback summary

- reinsert feedback to align path characterisations at receiver
- packets arrive at each router predicting downstream path
- arranged for dominant strategy of all parties to be honesty
- incremental deployment + upgrade incentive knob
- hangs new capabilities on ECN deployment, not just performance
- a simple idea for the Internet's accountability architecture



- democratises path information
 - either network or source can control (control requires timely information)
 - designed for tussle: preserves e2e principle, but endpoint control optional



(non-)issues with layering & tunnels

- general non-issue
 - RE flag shouldn't change once set by sender (or proxy)
 - policers merely read RE to compare with CE introduced so far
 - OK as long as CE represents congestion since same origin that set RE
- IP in IP tunnels
 - OK if tunnel entry copies RE and CE to outer header
 - but full functionality RFC3168 ECN tunnel resets CE in outer header
 - RFC3168 only said reset because security folks thought copy might leak info
 - concern has been resolved – updated IPSec RFC4301 (Dec 05) copies ECN at ingress
 - RFC3168 tunnelling section needs updating to reflect later security thinking and practice
- IP payload encryption (e.g. IPSec ESP)
 - non-issue – re-ECN designed to work only in network layer header
 - flow-ID obfuscation also non-issue – re-ECN only uses flow ID uniqueness, if at all
- layer 2 congestion notification (ATM, Frame, ... MPLS, 802.3ar)
 - non-issue given IP layer should accumulate CE from each 'L2 network' into ECN

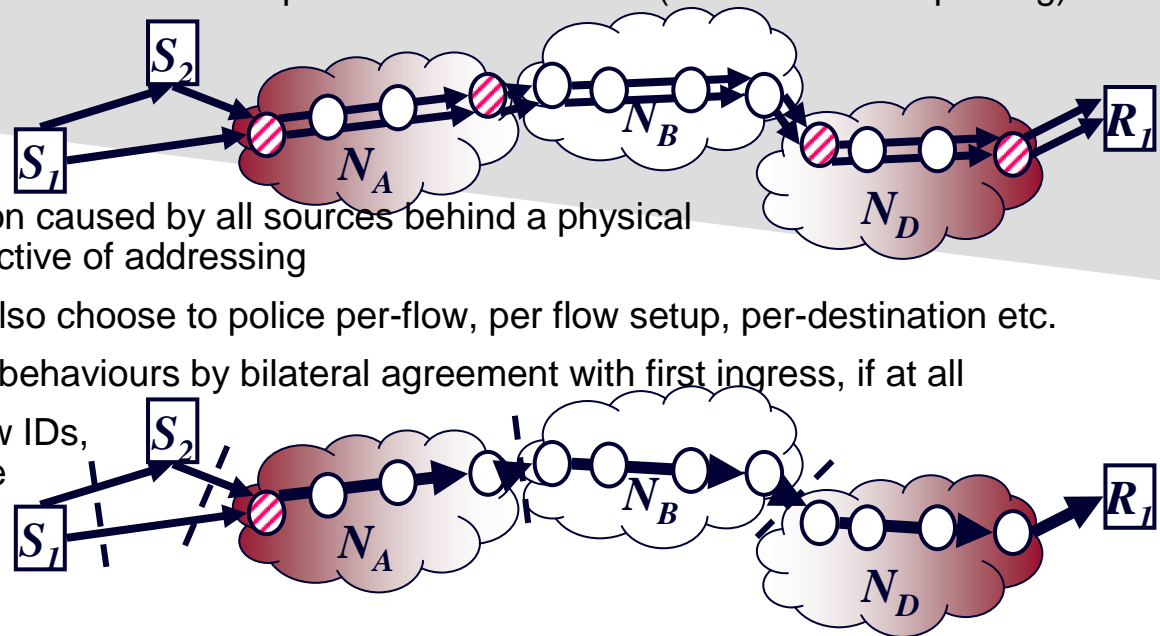
bottleneck policing harmful to evolvability

...and bypass-able anyway

- bottleneck policers: active research area since 1999
 - detect misbehaving flows causing 'unfair' share of congestion
 - located at each potentially congested routers
 - what right have these policers to assume a specific congestion response for a flow?
 - if they could police accurately, new congestion control evolution would require per-flow authorisation from all policers on the path (cf. IntServ)
 - malicious sources can bypass them by splitting flow IDs
 - even splitting flow across multiple intermediate hosts (or src address spoofing)

- re-ECN policing

- polices congestion caused by all sources behind a physical interface, irrespective of addressing
- within that, can also choose to police per-flow, per flow setup, per-destination etc.
- evolution of new behaviours by bilateral agreement with first ingress, if at all
- dropper uses flow IDs, but no advantage to split IDs



independence from identifiers

- controls congestion crossing any physical interface
 - user-network, network-network
 - congestion from network layer down to physical
 - not from a source address
- does have a dependency on source addresses
 - not to **identify** sources, merely to treat each flow separately
 - outstanding vulnerability
 - attacker spoofs another source's flow
 - deliberately brings down their joint average causing high drop

extended ECN codepoints: summary

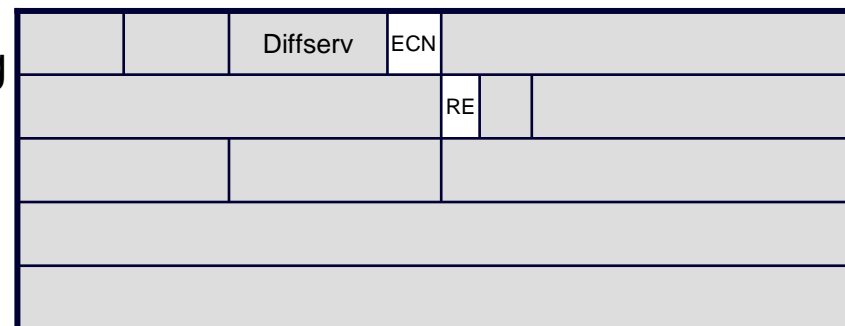
- extra semantics backward compatible with previous ECN codepoint semantics

ECN code-point	ECN [RFC3168] codepoint	RE flag	Extended ECN codepoint	re-ECN meaning	`worth`
00	not-ECT	0	Not-RECT	Not re-ECN capable transport	
		1	FNE	Feedback not established	+1
01	ECT(1)	0	Re-Echo	Re-echo congestion event (ECN nonce conflict)	+1
		1	RECT	Re-ECN capable transport	0
10	ECT(0)	0	---	'Legacy' ECN use	
		1	--CU--	Currently unused	
11	CE	0	CE(0)	Congestion experienced with Re-Echo	0
		1	CE(-1)	Congestion experienced	-1

re-ECN wire protocol in IPv4 (§3)

- propose Re-ECN Extension (RE) flag

- for IPv4: propose to use bit 48 (was reserved)
- set by sender, unchanged e2e

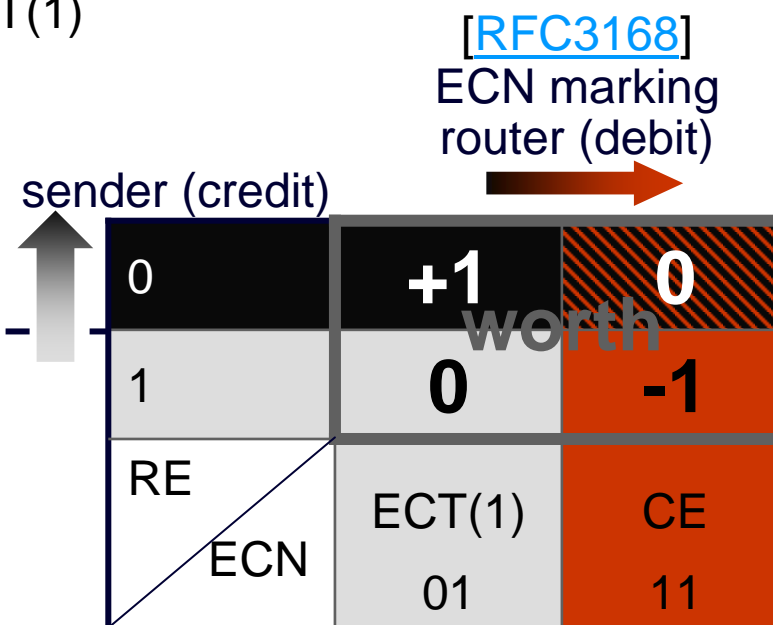


- once flow established
- sender re-inserts ECN feedback into forward data ("re-ECN") as follows
 - re-ECN sender always sets ECT(1)

- on every **congestion event** from transport (e.g. TCP)

sender blanks RE
else sets RE

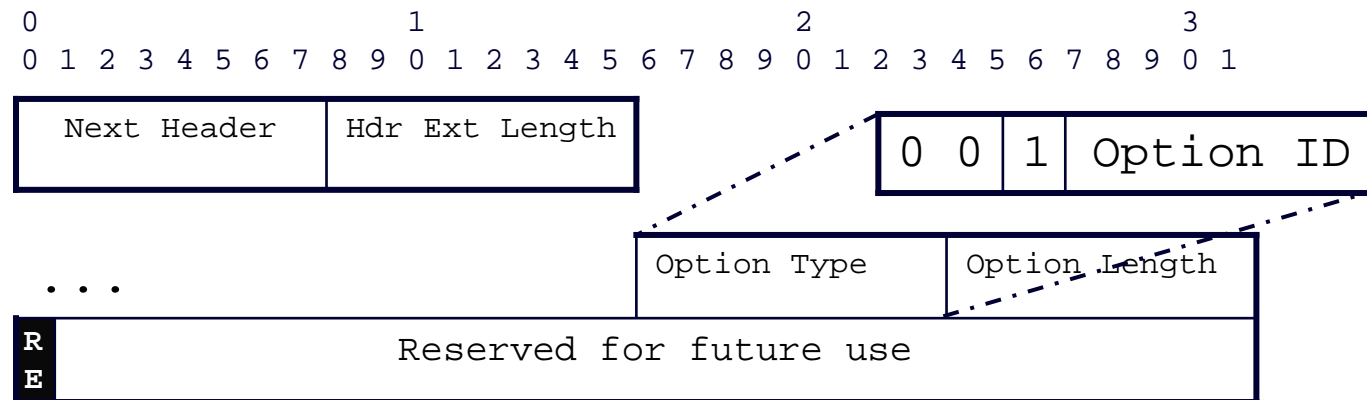
- conceptually, 'worth' of packet depends on 3 bit 'codepoint'
- aim for zero balance of worth in flow



IPv6 re-ECN protocol encoding

- IPv6 hop-by-hop options header extension

- new Congestion hop-by-hop option type



- action if unrecognized (AIU) = 00 'skip and continue'
 - changeable (C) flag = 1 'may change en route'
 - even tho RE flag shouldn't change en route (AH would just tell attackers which packets not to attack)
- seems wasteful for 1 bit, but we plan:
 - future hi-speed congestion control I-D using multi-bit congestion field
 - other congestion-related fields possible
 - e.g. to distinguish wireless loss and per-packet vs per-bit congestion

OPTIONAL router forwarding changes

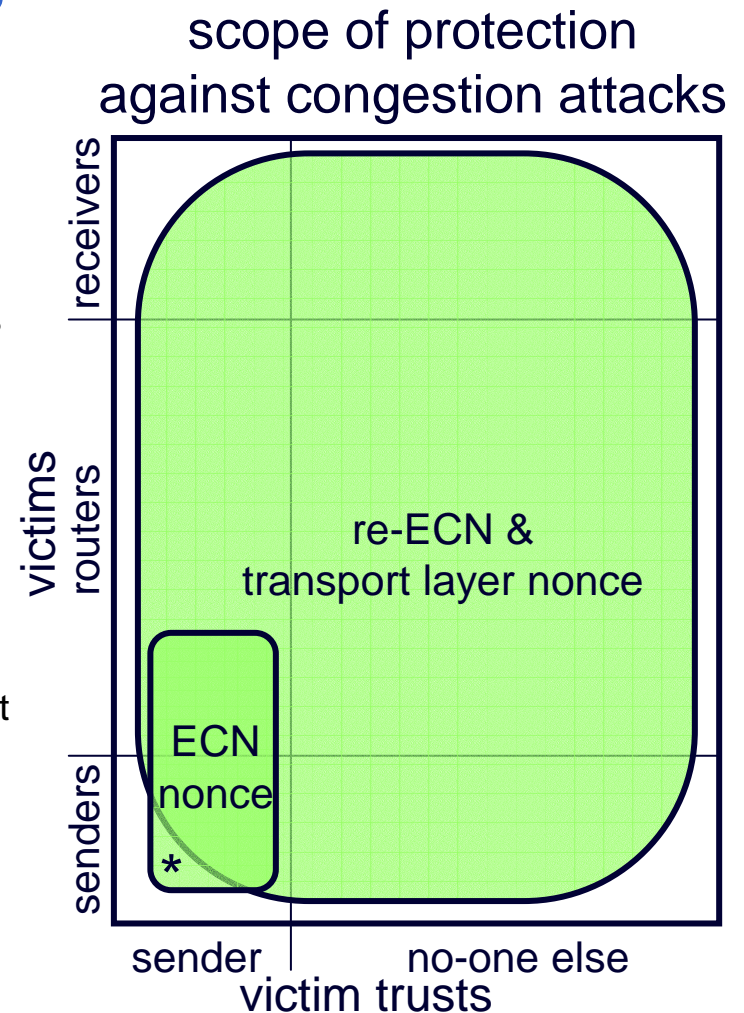
- preferential drop: improves robustness against DDoS
- **green** can be ECN marked rather than dropped (with caveat)

ECN code-point	ECN [RFC3168] codepoint	RE flag	Extended ECN codepoint	re-ECN meaning	`worth'	pref drop (1=drop 1 st)
00	not-ECT	0	Not-RECT	Not re-ECN capable transport		1
		1	FNE	Feedback not established	+1	3
01	ECT(1)	0	Re-Echo	Re-echo congestion event	+1	3
		1	RECT	Re-ECN capable transport	0	2
10	ECT(0)	0	---	'Legacy' ECN use		1
		1	--CU--	Currently unused		1
11	CE	0	CE(0)	CE with Re-Echo	0	2
		1	CE(-1)	Congestion experienced	-1	2

new appendix “Argument for holding back the ECN nonce” (§AI)

ECN nonce usefulness

- re-ECN and a transport layer nonce defend against wide range of attacks
 - ECN nonce defends against a small subset
 - and only one outside re-ECN’s range (*)
 - a sender that uses network ECN to allocate its own resources, can limit a lying receiver
 - sender can contain this attack without nonce
- IP header bits used to do this:
 - ECN nonce $\frac{1}{4}b$ (leaving last bit)
 - re-ECN $\frac{3}{8}b$ (using last bit)
- one common codepoint
 - re-ECN negotiates its use, but ECN nonce doesn’t
- propose to hold back ECN nonce
 - to see if we can find a coding to do both
 - to see if we can prevent (*) another way
 - develop a transport layer solution



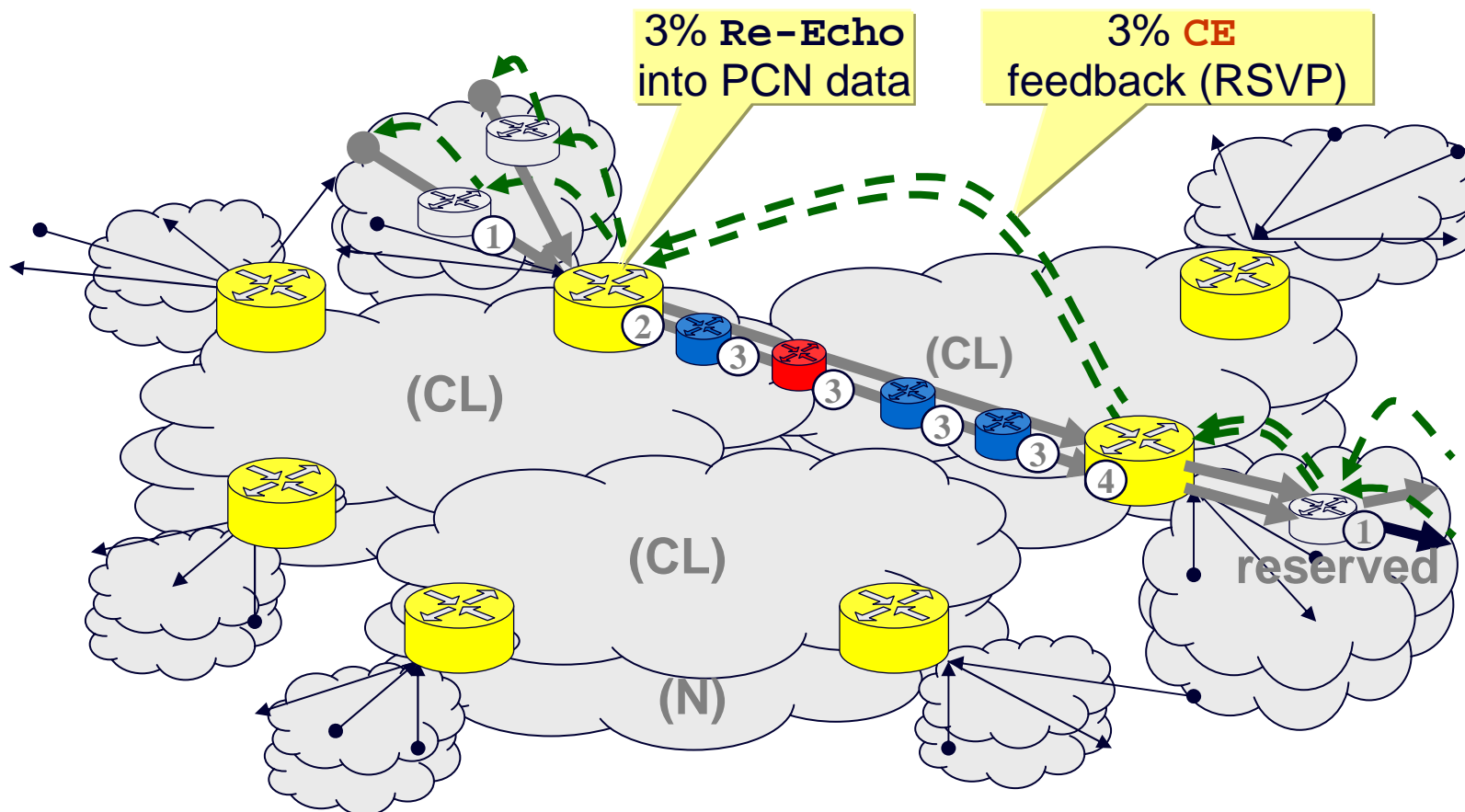
flow bootstrap

- at least one **green** packet(s) at start of flow or after >1sec idle
 - means “feedback not established”
 - ‘credit’ for safety due to lack of feedback
 - a **green** byte is ‘worth’ same as a **black** byte
- a different colour from black
 - distinguishes expected congestion based on experience from based on conservatism
 - gives deterministic flow state mgmt (policers, droppers, firewalls, servers)
 - rate limiting of state set-up
 - congestion control of memory exhaustion
- **green** also serves as state setup bit [Clark, Handley & Greenhalgh]
 - protocol-independent identification of flow state set-up
 - for servers, firewalls, tag switching, etc
 - don’t create state if not set
 - may drop packet if not set but matching state not found
 - firewalls can permit protocol evolution without knowing semantics
 - some validation of encrypted traffic, independent of transport
 - can limit outgoing rate of state setup
- to be precise **green** is ‘idempotent soft-state set-up codepoint’

guidelines for adding re-ECN to other transports

- main focus of <[draft-briscoe-tsvwg-re-ecn-tcp-03](#)>
 - IP (§5)
 - TCP (§4.1)
- added very brief sections giving guidelines for
 - DCCP (§4.2.3)
 - SCTP (§4.2.4)
 - spec would have to be a new I-D in each case
- focus of <[draft-briscoe-tsvwg-re-ecn-border-cheat-01](#)>
 - RSVP/NSIS transports ('re-PCN')
 - proposed technique to extend PCN-based admission control
 - Internet wide (edge-edge) – many untrusting domains
- our current focus
 - controlling fairness between current transports & hi-speed congestion control

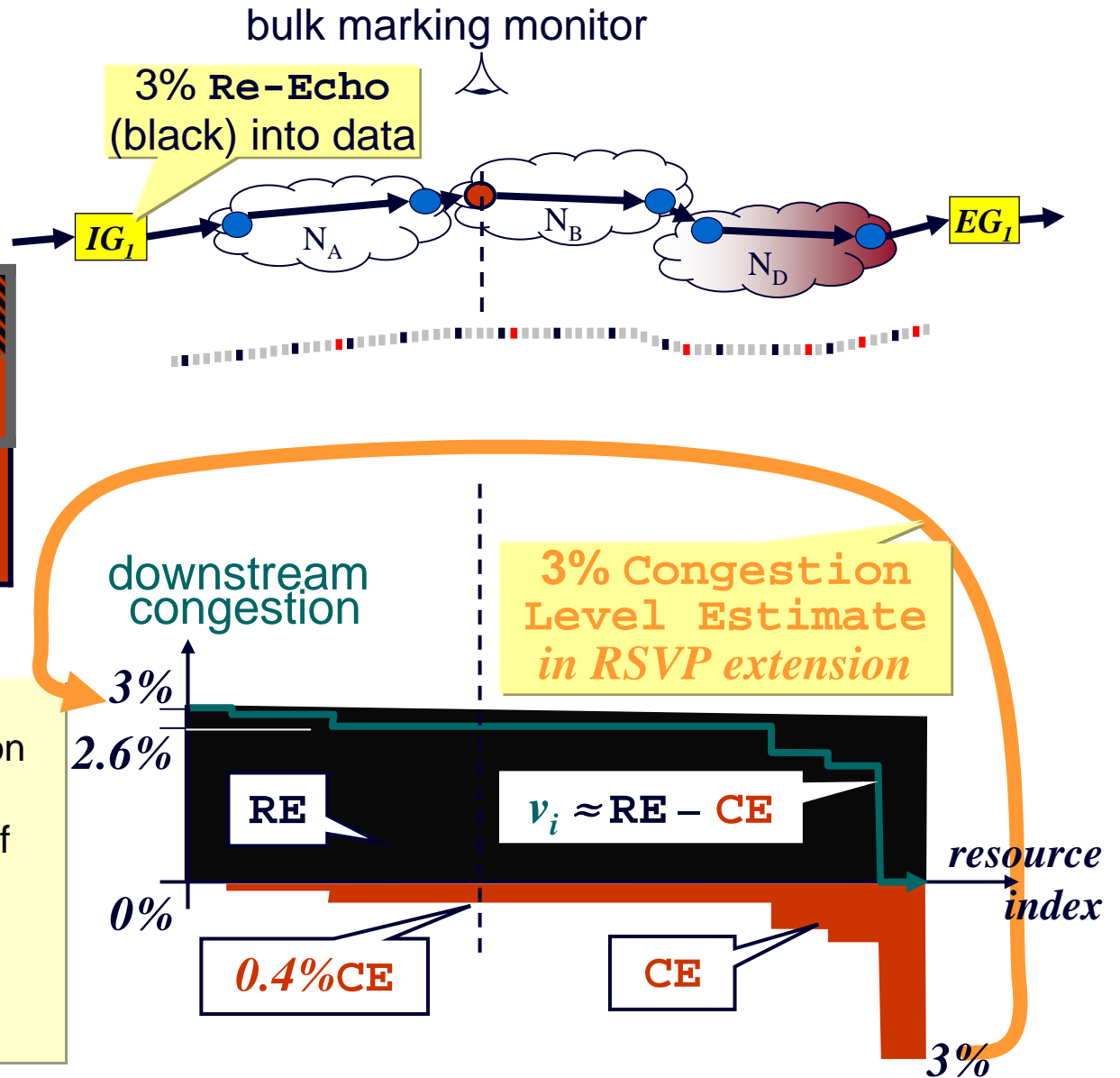
border anti-cheating solution



re-PCN

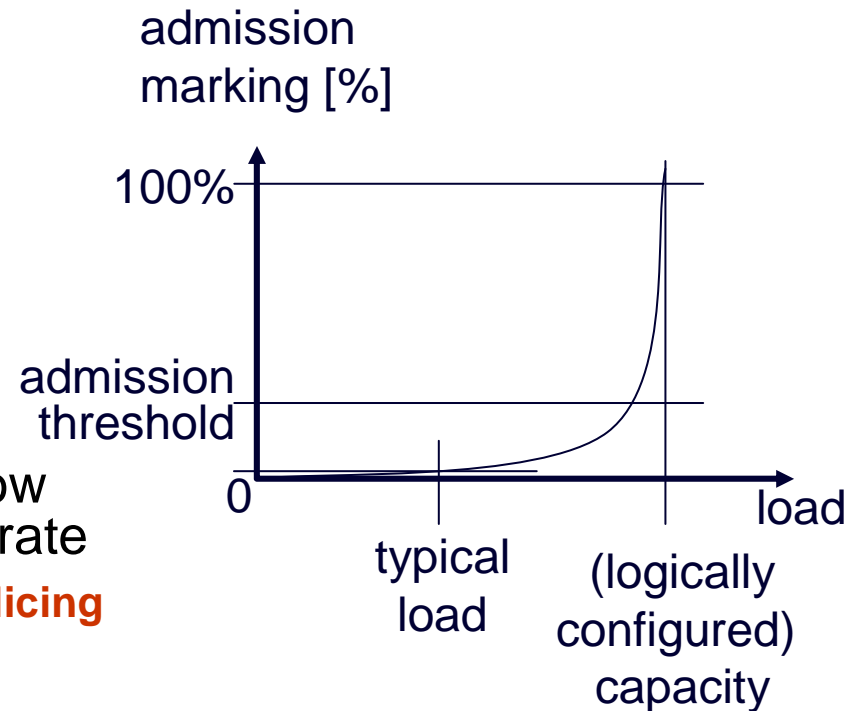
0	+1	0
1	0	-1
RE ECN	ECT(1) 01	CE 11

- ingress gateway blanks **RE**, in same proportion as fraction of **CE** arriving at egress
- at any point on path, bulk diff betw fractions of **RE** & **CE** is downstream congestion
- routers unchanged



solution rationale

- <0.01% packet marking at typical load
 - addition of any flow makes little difference to marking
- penalties to ingress of each flow appear proportionate to its bit rate
 - **emulates border flow rate policing**
- as load approaches capacity
 - penalties become unbearably high (~1000x typical)
 - insensitive to exact configuration of admission threshold
 - **emulates border admission control**
- neither is a perfect emulation
 - but should lead to the desired behaviour
 - fail-safes if networks behave irrationally (e.g. config errors) – see draft

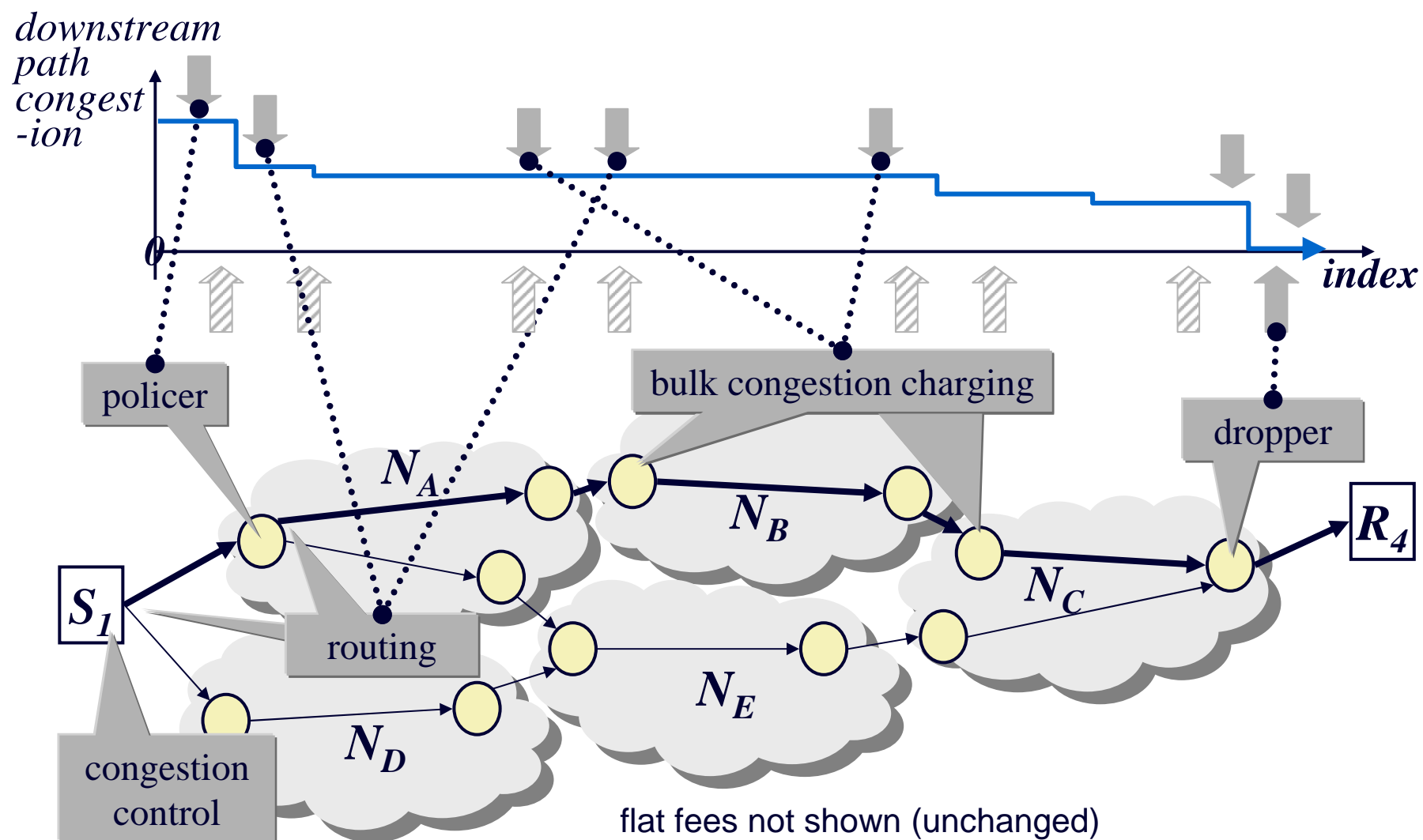


differential quality of service (QoS) control without all the complicated stuff

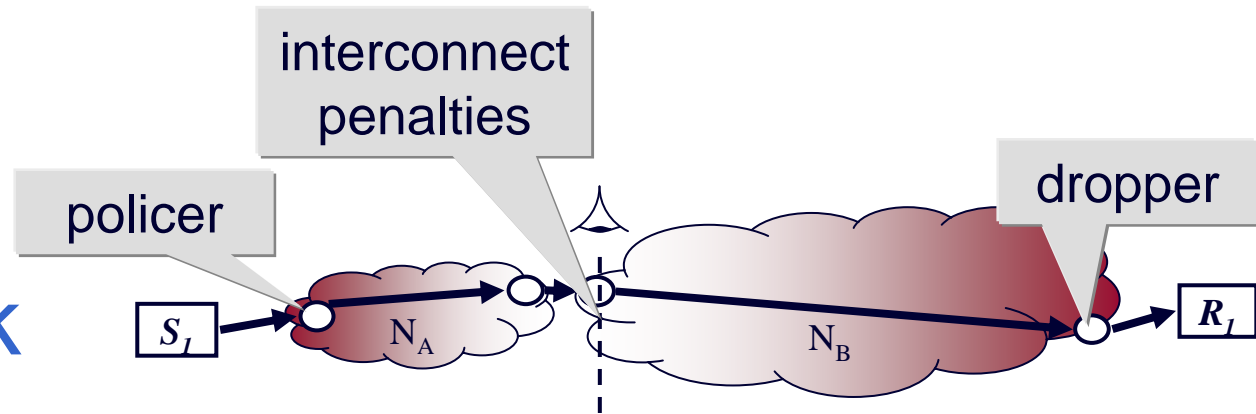
- QoS only relevant when there's a risk of congestion
- enforcing congestion control is equivalent to QoS (and to paying for it)
 - allowing one app's rate to slow down less than others in response to incipient congestion (ie. still low delay)
 - is equivalent to giving scheduling priority on routers*
- even if user pays a flat monthly fee
 - better QoS for some apps leaves less congestion 'quota' for rest
- purely by local (sender↔ingress) arrangement
 - no authorisation on any other network elements (equal marking)
- other networks reimbursed automagically
 - by inter-domain congestion pricing (SLA model also possible)
- incredible simplification of mechanisms for QoS control & mgmt
 - and, unlike other QoS mechanisms
 - it also prevents users 'stealing' QoS at everyone else's expense

* except within a round trip time – implies two priority classes would be sufficient
(can also determine relative congestion marking rates of each class using economics)

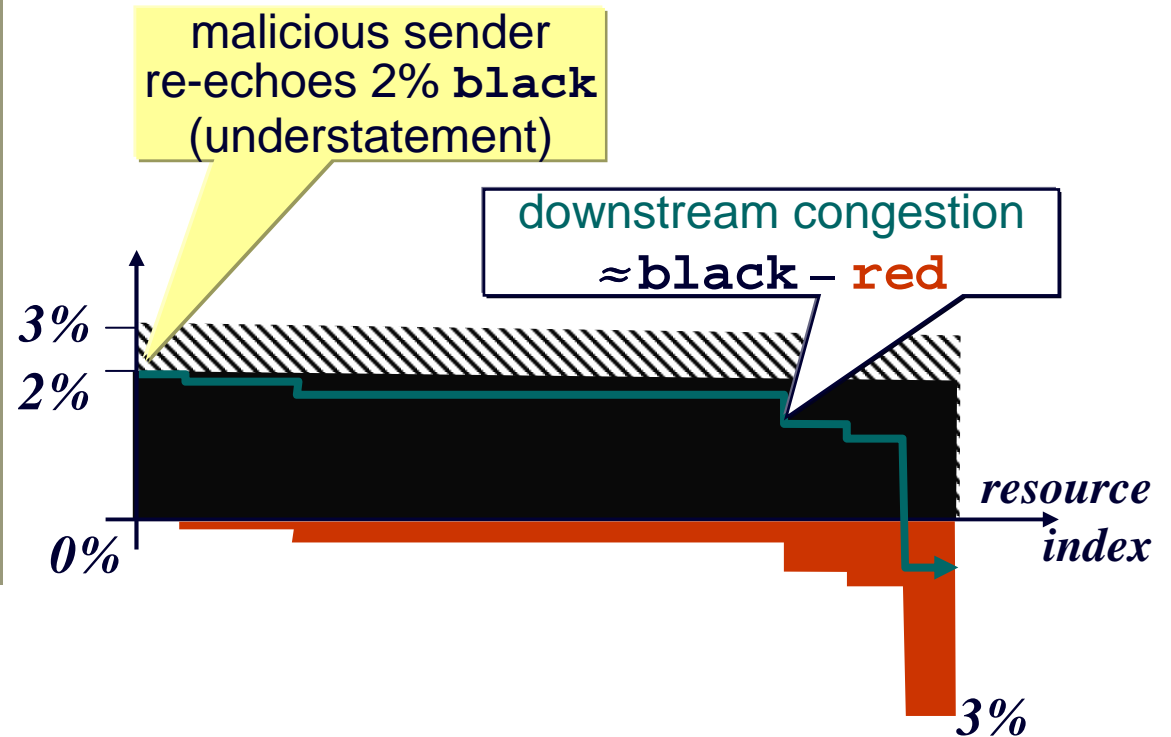
incentive framework



incentive framework

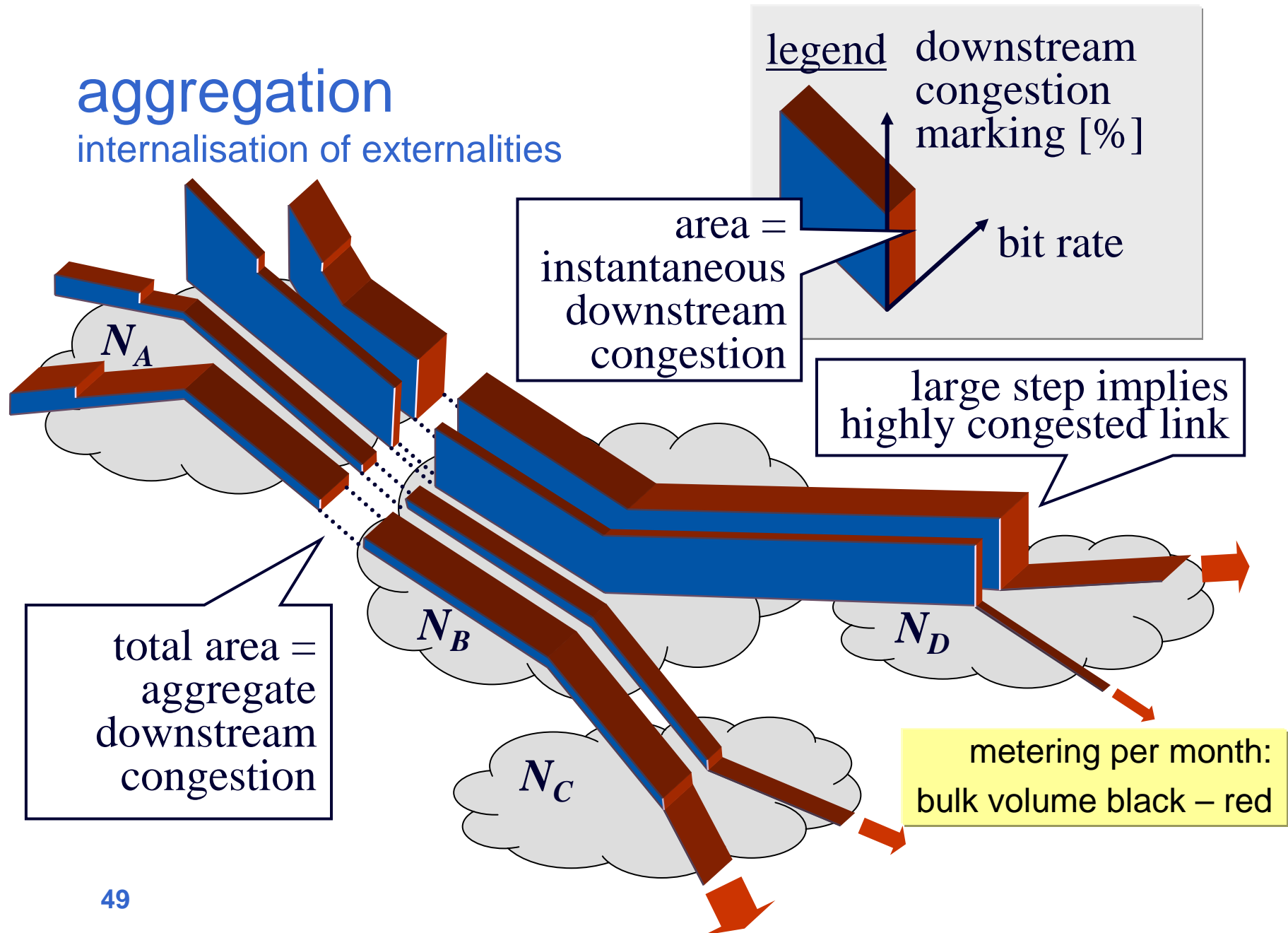


- packets carry view of downstream path congestion to each router
- using path congestion declared by sender
 - can police rate response
 - or enforce congestion quotas
- won't sender or rcvr just understate congestion?
 - egress drops negative balance (next slide)



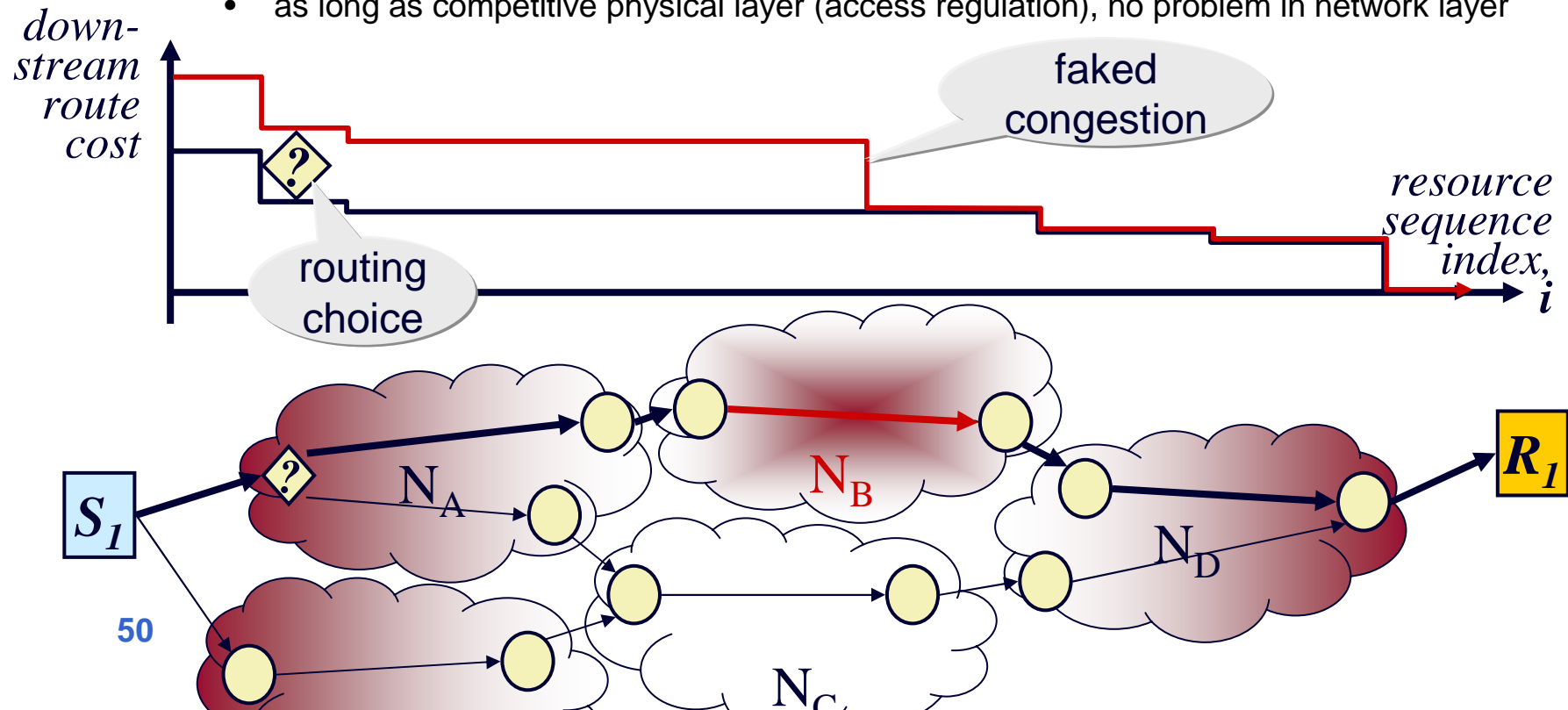
aggregation

internalisation of externalities



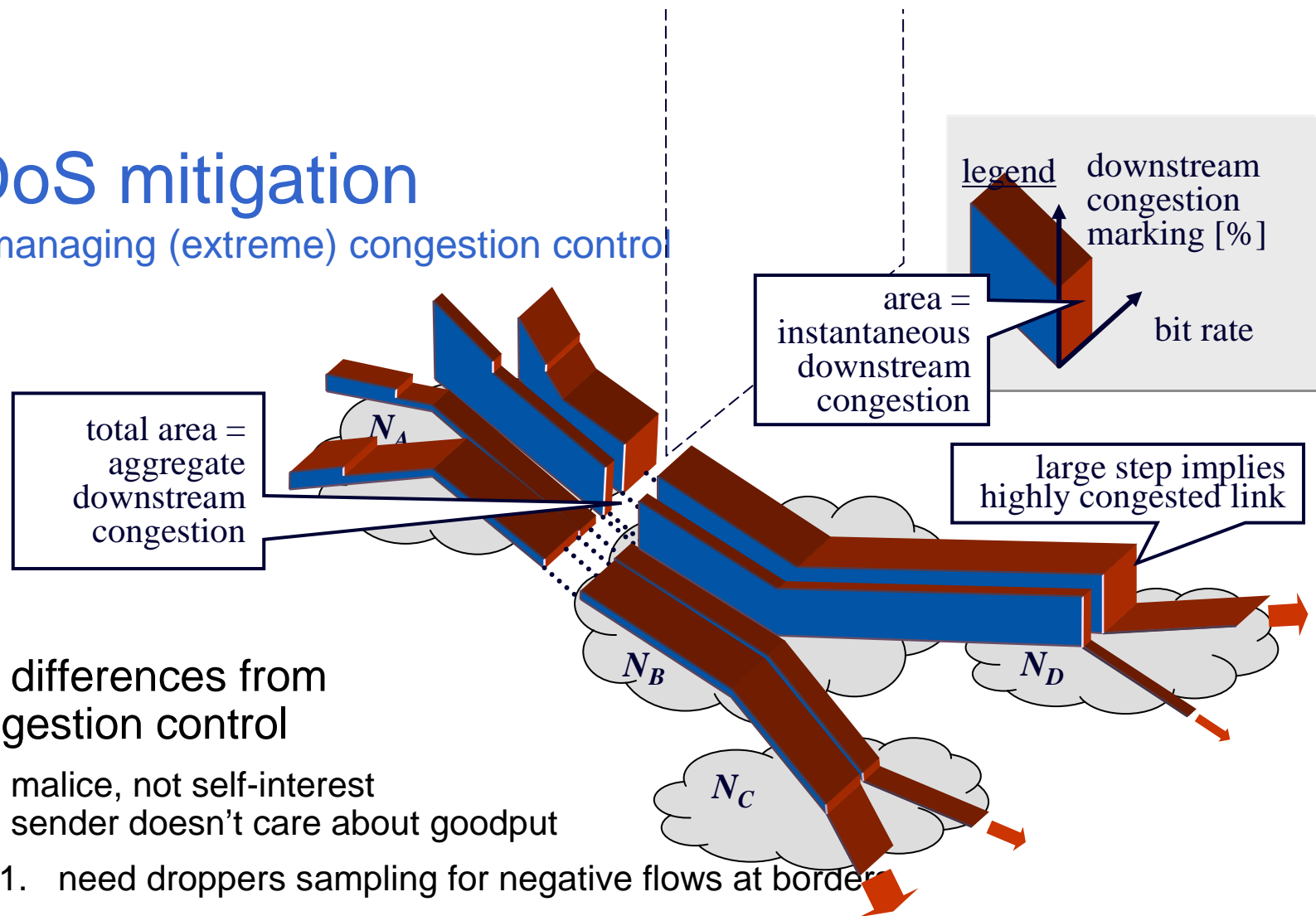
congestion competition – inter-domain routing

- if congestion \rightarrow profit for a network, why not fake it?
 - upstream networks will route round more highly congested paths
 - N_A can see relative costs of paths to R_1 thru N_B & N_C
- the issue of monopoly paths
 - incentivise new provision
 - as long as competitive physical layer (access regulation), no problem in network layer



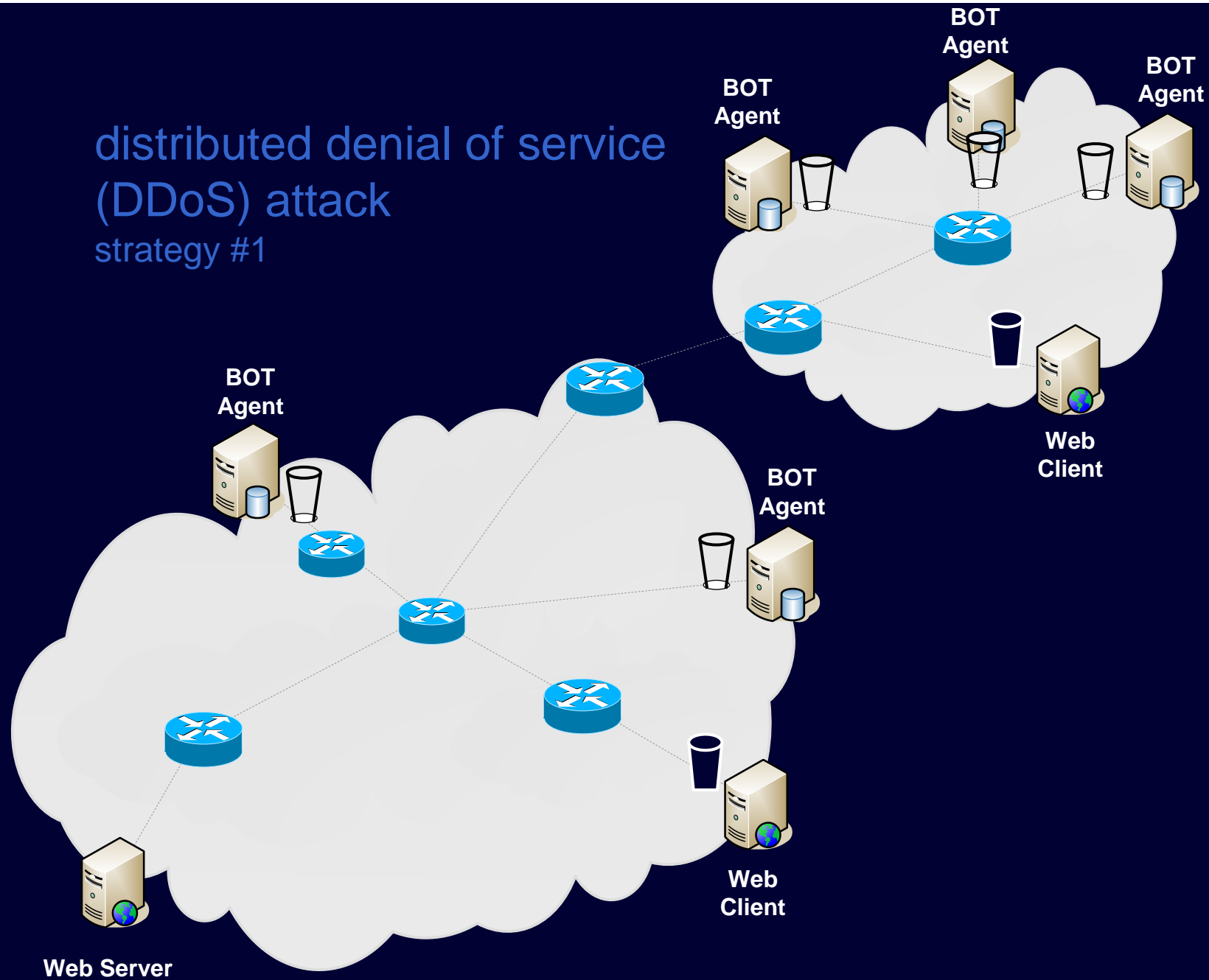
DDoS mitigation

just managing (extreme) congestion control



- two differences from congestion control
 - malice, not self-interest
sender doesn't care about goodput
 1. need droppers sampling for negative flows at border
 - pushes beyond incipient congestion into heavy loss
 2. need preferential drop on routers
- provides incentives to deploy complementary DDoS solutions

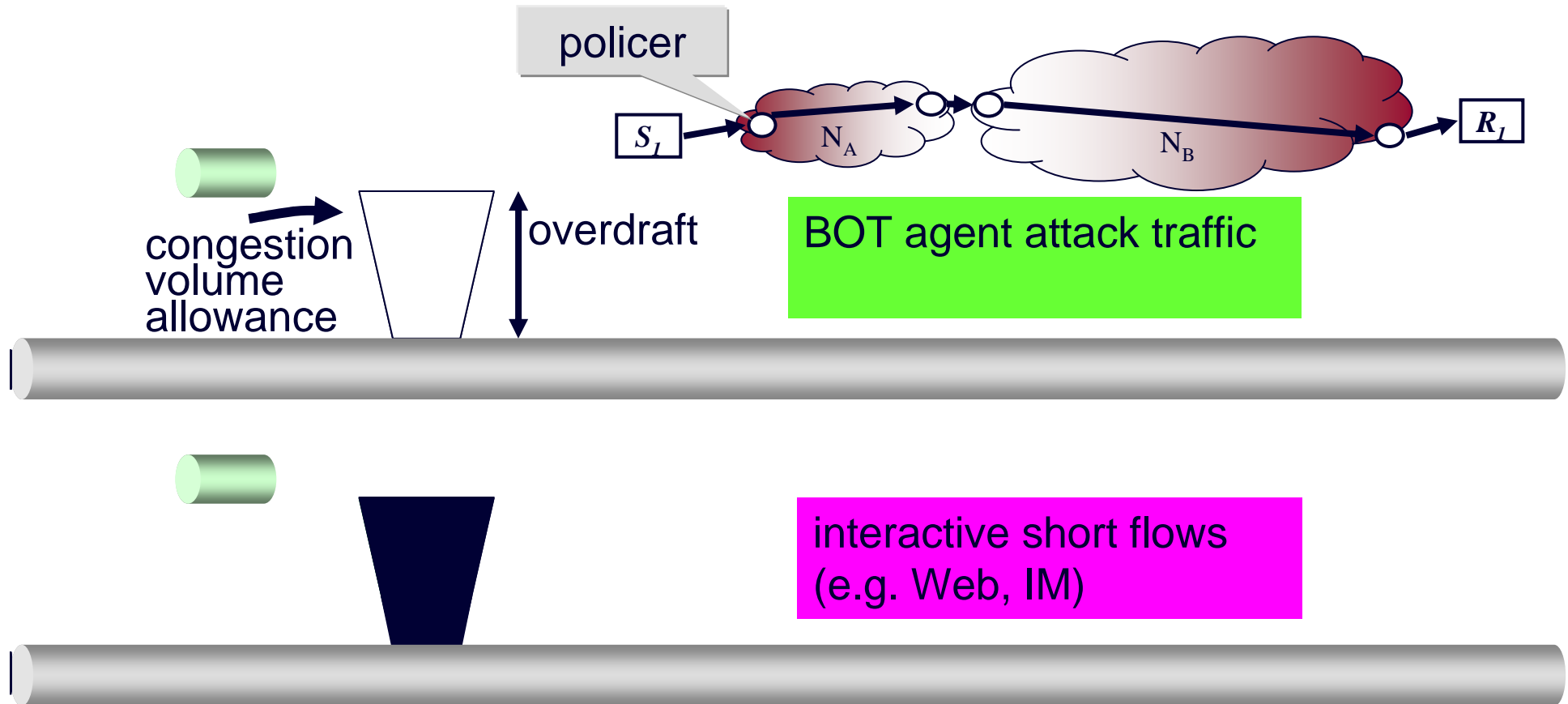
distributed denial of service (DDoS) attack strategy #1



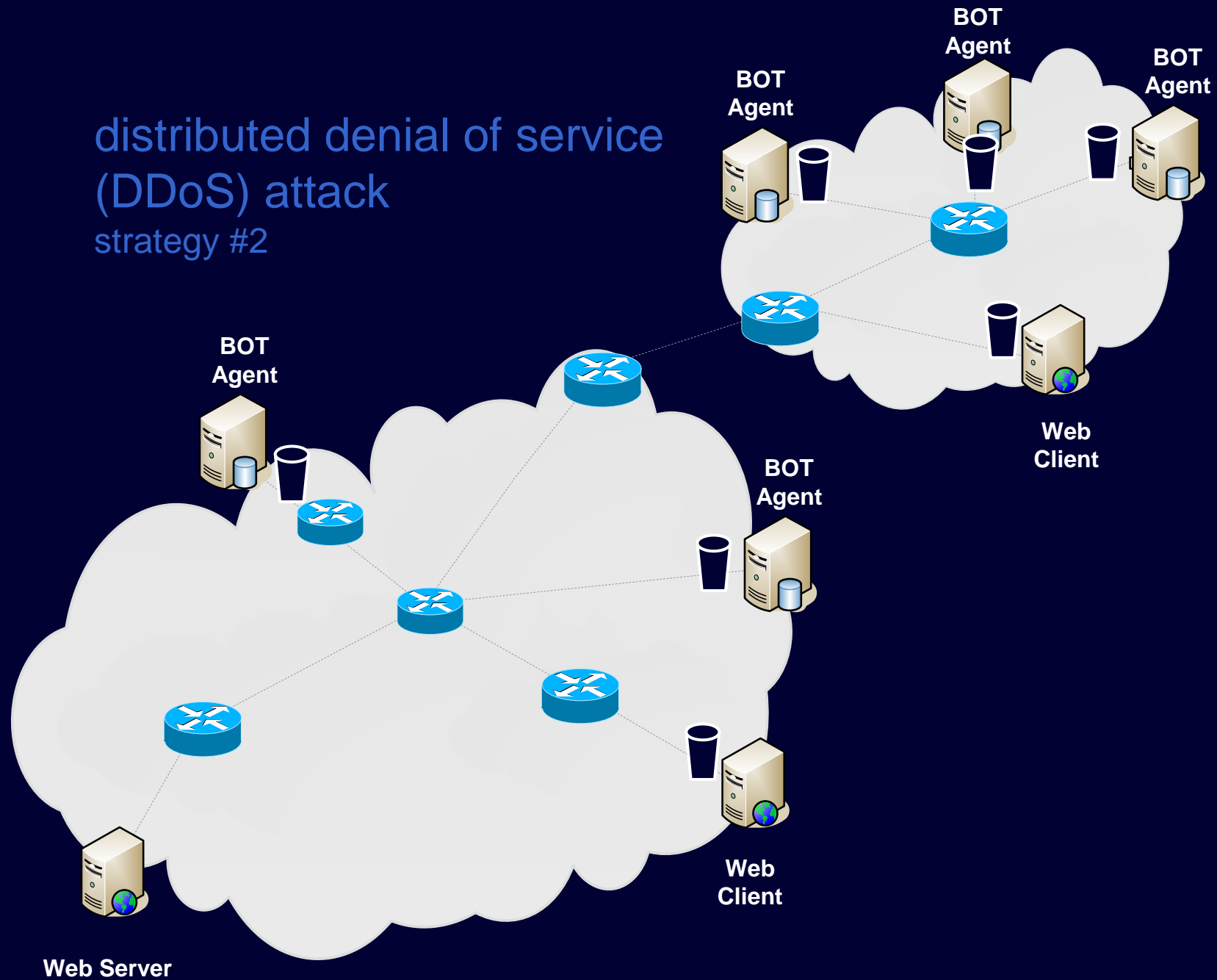
animation requires Office XP or equivalent

per-user congestion policer

DDoS attack strategy #1



distributed denial of service (DDoS) attack strategy #2



animation requires Office XP or equivalent

attacks on re-ECN & fixes

0	+1	0
1	0	-1
RE ECN	ECT(1) 01	CE 11

- recap: why two codepoints worth 0?
 - when no congestion send neutral (0)
 - packet marked 'cancelled' if network happens to mark a packet (-1) which the sender used to re-echo congestion (+1); $+1 - 1 = 0$
 - in draft 00, congestion marking of +1 packet turned it to -1 not 0, but networks could cheat by focusing marking on +1 (see §B)
- but now can't attacker just send cancelled packets?
 - immune from congestion marking
 - simple fix: policer counts cancelled with +1 towards *path* congestion
 - should have specified this anyway, as both represent path congestion
 - also check proportion of cancelled to +1 packets same as -1 to neutral
- set of attacks using persistently negative dummy traffic flows
 - see next presentation for border policing fix
- one remaining known vulnerability if attacker can spoof another flow ID
 - known since early on – plan to focus effort on fixing this next

dummy traffic attacks on re-ECN

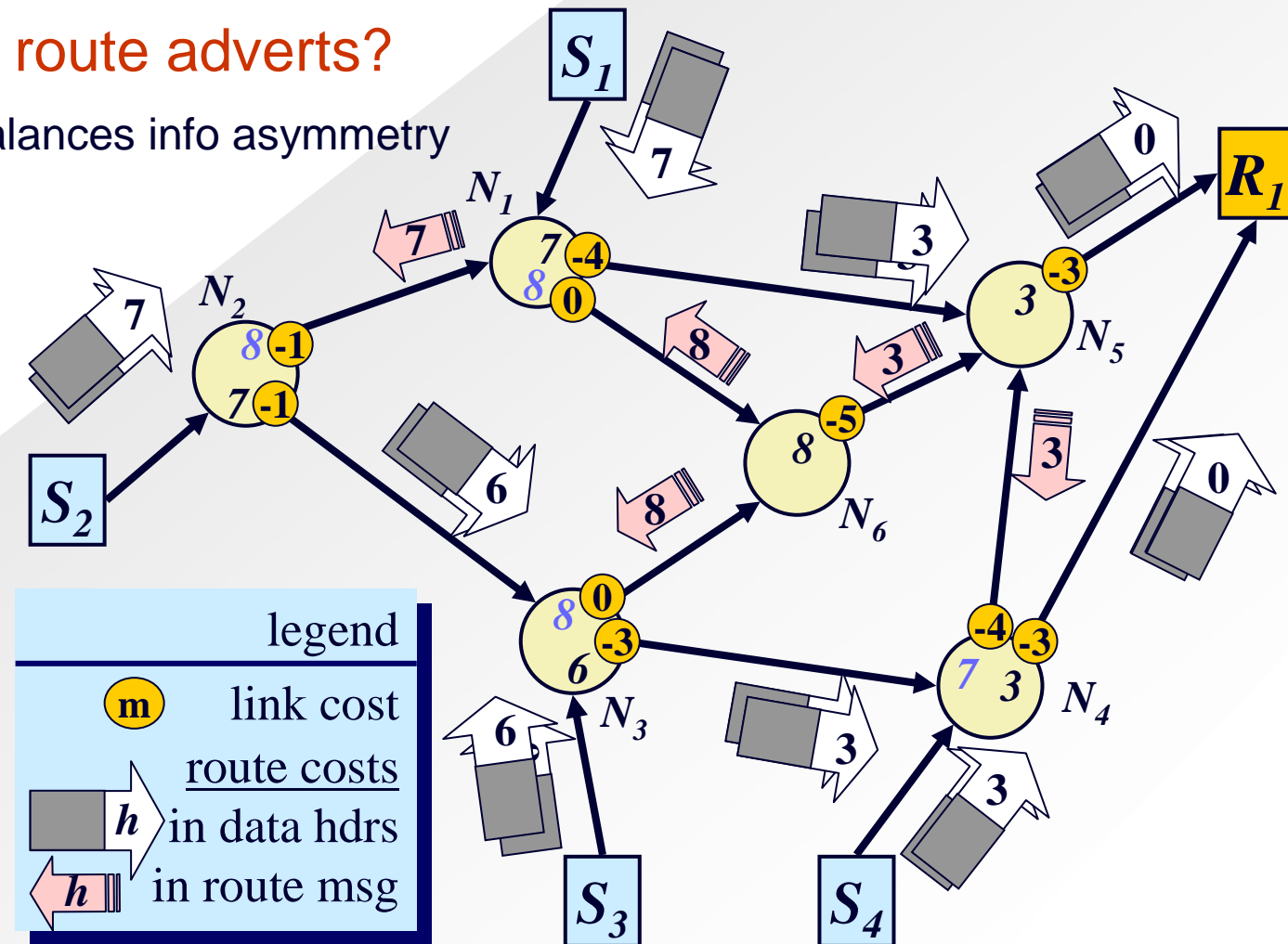
- sanctions against persistently negative flows may not discourage dummy traffic
- various attacks ([Salvatori, Bauer] see draft), eg.
 - a network sends negative dummy traffic with just enough TTL to cross border [Salvatori]
 - offsets penalties from other positive traffic
- fix is to estimate contribution from negative flows crossing border by sampling
 - inflate penalties accordingly – removes attack motivations
 - see draft for details and example algorithm in appendix

re-ECN security considerations (§10) and incentive framework limitations (§6.3)

- egress dropper
 - robust against attack that plays-off against ingress policing
 - robust against state exhaustion attacks (by design of **green**)
 - write-up of state aggregation implementation TBA
 - believe new protocol allows dropper to be robust against dynamic attacks
- collateral damage attack still possible → next slide
- re-ECN deliberately designed not to rely on crypto

load balanced routing support?

- automate inter-domain traffic engineering (damped)?
- **validate route adverts?**
 - re-balances info asymmetry



BT IPR related to [draft-briscoe-tsvwg-re-ecn-tcp-00.txt](https://datatracker.ietf.org/public/ipr_detail_show.cgi?&ipr_id=651)

- See IPR declaration at https://datatracker.ietf.org/public/ipr_detail_show.cgi?&ipr_id=651 which overrides this slide if there is any conflict
- 1) WO 2005/096566 30 Mar 2004 published
- 2) WO 2005/096567 30 Mar 2004 published
- 3) PCT/GB 2005/001737 07 May 2004 published
- 4) GB 0501945.0 (EP 05355137.1) 31 Jan 2005 published
- 5) GB 0502483.1 (EP 05255164.5) 07 Feb 2005 published
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more info...

- Fixing mindset on fairness
 - [Flow Rate Fairness: Dismantling a Religion](#) IETF Internet draft (Mar 2007)
- Overall re-feedback idea, intention, policing, QoS, load balancing etc
 - [Policing Congestion Response in an Inter-Network Using Re-Feedback](#) (SIGCOMM'05 – mechanism outdated)
- Protocol Spec and rationale
 - [Re-ECN: Adding Accountability for Causing Congestion to TCP/IP](#) IETF Internet Draft (Oct 2006)
- Using re-ECN with pre-congestion notification (PCN)
 - [Emulating Border Flow Policing using Re-ECN on Bulk Data](#) IETF Internet draft (Jun 2006)
- Relation between re-ECN and inelastic QoS
 - [Commercial Models for IP Quality of Service Interconnect](#) BT Technology Journal (Apr 2005)
- Mitigating DDoS with re-ECN
 - [Using Self-interest to Prevent Malice; Fixing the Denial of Service Flaw of the Internet](#) Workshop on the Economics of Securing the Information Infrastructure (Oct 2006)
- more related papers and all the above:
<<http://www.cs.ucl.ac.uk/staff/B.Briscoe/projects/refb/>>