# PI<sup>2</sup>: A Linearized AQM for both Classic and Scalable TCP

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# Use DCTCP on the Internet ?

- consistently low queuing delay
- full link utilization with very small queues
- ☺ very low loss
- more stable throughput between competing flows
- scalable with higher link rates
- available in Windows 10 and Linux 3.18
- not yet optimal (for high RTTs, ...)

# Can we use DCTCP on the Internet ?

Unfortunately (currently) not:

Starves the classic TCP-friendly flows

😕 keeps big tail drop queues full

😕 needs ECN, so high loss (or fallback to Reno)

Only used where everything can change at once

→ gradual & safe migration strategy

## Challenges

How to:

#### make DCTCP and TCP-Reno rate compatible $\rightarrow$ this paper: PI<sup>2</sup> AQM

preserve low latency for DCTCP
→ next papers: DualQ-PI<sup>2</sup> & TCP-Prague
→ IETF: L4S BoF successful, drafts in tsvwg

#### AQM for multiple congestion controls



F, G adapt the probability to align the steady state congestion window of both CCs

#### Equal congestion window



#### AQM for DCTCP and Cubic



Implemented as a Linux tc qdisc: <u>https://github.com/olgabo/dualpi2</u> Evaluated on a real testbed

#### AQMs for steady state test results



### Equal rate at different RTTs



# Equal rate with different flow nbrs

PIEECN-Cubic(A)/Cubic(B) ratioPIEDCTCP(A)/Cubic(B) ratioPI2DCTCP(A)/Cubic(B) ratio



Nr of flows for each cc(A-B)

✓ Equal window for steady state

- ? Dynamic behavior
- ? Stability PI

# PI-AQM recap

Every  $T_{update}$  interval do:

 $\Delta p = \alpha^*(error) + \beta^*(queue change)$ 



# Choosing $\alpha$ and $\beta$

Larger  $\alpha$  and  $\beta$  values give faster response

Stability analysis: stable if gain margin > 0



Gain margin evolves diagonally with  $p \rightarrow$  problem!

#### PIE solution: $\alpha$ and $\beta$ tuning

Adapt (tune)  $\alpha$  and  $\beta$  based on p



# PIE solution: $\alpha$ and $\beta$ tuning

Tune  $\alpha$  and  $\beta$  based on previous p:



Works well. Tuning is required improvement ! Curing the symptoms



# PI<sup>2</sup> solution: remove the $\sqrt{}$

Stability models used for:

TCP Reno on PI:  $\frac{dW(t)}{dt} = \frac{1}{R(t)} - 0.5 \frac{W(t)W(t - R(t))}{R(t - R(t))} p(t - R(t))$ TCP Reno on PI2:  $\frac{dW(t)}{dt} = \frac{1}{R(t)} - 0.5 \frac{W(t)W(t - R(t))}{R(t - R(t))} (p'(t - R(t)))^2$ 30 reno pie 20 Gain Margin [dB] reno pi2 scal pi 10 0 -10 unstable -20 -30 10 100 0.1 1 p' [%]

#### Effect on not squaring PI for Reno



## PI<sup>2</sup> similar to PIE for Reno



Time [sec]

# PI(<sup>2</sup>) controls DCTCP



Time [sec]

# Conclusions

PI<sup>2</sup> is simpler than PIE, performs not worse and supports scalable CCs (without the square)

PI controls natively scalable CCs, use adaption function to convert any CC to a scalable

Future work:

Single Q deployment is not recommended for low latency

- $\rightarrow$  DualQ to preserve low latency
- $\rightarrow$  TCP-Prague to improve DCTCP for Internet

#### Questions?

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https://github.com/olgabo/dualpi2 http://riteproject.eu/dctth

# Backup

## DCTCP recap

TCP (Reno) ←→ DCTCP

#### Response to congestion in sender

Half the congestion window

On average half a packet per ECN mark

→ React to level of congestion

ECN feedback in receiver

Echo once per RTT

Echo every mark / non-mark

→ accurate ECN feedback

#### ECN marking in network

Smooth and delay a drop or mark to allow bursts

Don't smooth or delay queue size→ immediate ECN marking

# DCTCP recap

#### **RED** for Reno

